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THE DIALECTICA INTERPRETATION OF FIRST-ORDER CLASSICAL AFFINE LOGIC

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ABSTRACT. We give a Dialectica-style interpretation of first-order classical affine logic. By moving to a contraction-free logic, the translation (a.k.a. D-translation) of a first-order formula into a higher-type $\exists\forall$-formula can be made symmetric with respect to duality, including exponentials. It turned out that the propositional part of our D-translation uses the same construction as de Paiva's dialectica category $\mathbb{GC}$ and we show how our D-translation extends $\mathbb{GC}$ to the first-order setting in terms of an indexed category. Furthermore the combination of Girard's ?!-translation and our D-translation results in the essentially equivalent $\exists\forall$-formulas as the double-negation translation and Gödel's original D-translation.

1. Introduction

Gödel's Dialectica interpretation is a functional interpretation of Peano Arithmetic (PA) [Troelstra 1973, Hindley and Seldin 1986, Gödel 1990, Avigad and Feferman 1998]. He translated a formula of PA into a formula of the form $\exists u \forall x\, \alpha(u, x)$, where u and x are sequences of higher-type variables and α is a formula of the quantifier free system T of higher-type functionals. He then showed that for every provable formula of PA one can find higher-type terms $\boldsymbol{u}$ such that $\alpha(\boldsymbol{u}, x)$ is provable in T. The consistency of PA is thus reduced to the consistency of T, which Gödel thought is intuitively justifiable.

The translation consists of two parts. The first part is the Gödel-Gentzen double-negation translation ($\neg\neg$-translation) of classical logic into intuitionistic logic. The second part is the so-called Dialectica translation (D-translation) of the first-order intuitionistic arithmetic (HA) into formulas of the form $\exists u \forall x\, \alpha(u, x)$. For the purpose of this paper we focus only on the first-order logic, neglecting the arithmetic.

The formulas $\exists u \forall x\, \alpha(u, x)$ obtained by Gödel's D-translation easily become very much complicated. We analyze his D-translation and find that contraction is mainly responsible for this. We hence move to a contraction-free logic, specifically affine logic, although part of our work is carried out in linear logic regarded as a subsystem of affine logic. The D-translation can then be further decomposed and made symmetric with respect to duality, including exponentials.

It turned out that the propositional part of our D-translation uses the same con-

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struction as de Paiva's dialectica category $\mathbb{GC}$ [de Paiva 1991, de Paiva 2006], which is a categorical model of propositional classical linear logic. This leads us to the idea to achieve a categorical model of first-order classical linear logic. We indeed show how our D-translation extends $\mathbb{GC}$ to the first-order setting in terms of an indexed category.

Furthermore the combination of Girard's ?!-translation [Girard 1987] and our D-translation results in the essentially equivalent formulas, in the sense we will define later, as the $\neg\neg$-translation and Gödel's original D-translation with respect to formulas of classical logic. We may hence say that our D-translation is a refinement of Gödel's original one. We have even better correspondence if we consider Shoenfield's version of D-translation [Shoenfield 1967] instead of Gödel's.

Since our D-translation gives the essentially equivalent formulas as Gödel's original with respect to classical formulas it would not be fair to say that ours simplifies Gödel's in general. We would rather claim only that our D-translation gives us a more fine-grained control in the construction of witness terms and enables us to find simpler terms when the use of contraction is limited.

The symmetry of our D-translation suggests that it can be recapitulated in terms of games. We are hoping to be able to analyze the computational meaning of the Dialectica interpretation further in this direction.

2. Background and motivation

2.1. Gödel's Dialectica interpretation. Gödel's Dialectica interpretation (D-interpretation) is intended to be a technique to prove the consistency of *PA* [Troelstra 1973, Hindley and Seldin 1986, Gödel 1990, Avigad and Feferman 1998]. Gödel's idea is the use of higher-type functionals to reduce the complexity of quantifier alternations. Any formula of *PA* is first translated ($\neg\neg$-translation) into a formula of *HA*, and then further translated (D-translation) into a formula of the form

$$\exists u_1 \cdots \exists u_m \forall x_1 \cdots \forall x_n \, \alpha(u_1, \ldots, u_m, x_1, \ldots, x_n, z_1, \ldots, z_l)$$

where $\alpha(u_1, \ldots u_m, x_1, \ldots, x_n, z_1, \ldots, z_l)$ is a formula of the quantifier-free system T whose free variables are among possibly higher-type variables $u_1, \ldots, u_m, x_1, \ldots, x_n$ and number variables $z_1, \ldots, z_l$. The propositional connectives in $\alpha(u_1, \ldots, u_m, x_1, \ldots, x_n, z_1, \ldots, z_l)$ can be taken as *classical* ones since only decidable predicates are considered.

Hereafter we use a single letter v for the list of variables $v_1, \ldots v_m$, and simply write $\alpha(u, x, z)$ for the formula $\alpha(u_1, \ldots u_m, x_1, \ldots, x_n, z_1, \ldots, z_l)$. The list v may be empty. Furthermore we let vy stand for the list of application terms

$$v_1 y_1 \cdots y_n, \; v_2 y_1 \cdots y_n, \ldots, \; v_m y_1 \cdots y_n$$

while vy is the empty list if v is. The variables u and x in $\exists u \forall x \, \alpha(u, x, z)$ will be called *positive* and *negative*, respectively. We often suppress the free variables z and write simply $\exists u \forall x \, \alpha(u, x)$ for readability.

2.2. Gödel-Gentzen $\neg\neg$-translation. For a formula ϕ of *PA*, its $\neg\neg$-translation ϕ^N into *HA* is inductively defined as follows.

1. If ϕ is atomic then $\phi^N \equiv \neg\neg\phi$.
2. $(\phi \wedge \psi)^N \equiv \phi^N \wedge \psi^N$.
3. $(\phi \vee \psi)^N \equiv \neg\left(\neg\phi^N \wedge \neg\psi^N\right)$.
4. $(\phi \supset \psi)^N \equiv \phi^N \supset \psi^N$.
5. $(\forall x\, \phi)^N \equiv \forall x\, \phi^N$.
6. $(\exists x\, \phi)^N \equiv \neg\forall x\, \neg\phi^N$.

The negation $\neg\phi$ is regarded as the abbreviation of $\phi \supset \bot$ in both *PA* and *HA*.

2.3. Theorem. [Gödel-Gentzen] *If ϕ is provable from $\psi_1, \psi_2, \dots, \psi_n$ in PA then ϕ^D is provable from ${\psi_1}^D, {\psi_2}^D, \dots {\psi_n}^D$ in HA.*

2.4. Gödel's D-translation. For a formula ϕ of *HA*, its D-translation ϕ^D is inductively defined together with ϕ_D such that $\phi^D \equiv \exists u \forall x\, \phi_D$. For formulas ϕ and ψ which appear in the inductive clauses we rename the bound variables if necessary so that we have

$$\phi^D \equiv \exists u \forall x\, \phi_D(u, x), \quad \psi^D \equiv \exists v \forall y\, \psi_D(v, y)$$

where there is no overlapping of variables among u, x, v, y and free variables in ϕ^D or ψ^D.

1. If ϕ is atomic then $\phi^D \equiv \phi$.
2. $(\phi \wedge \psi)^D \equiv \exists uv \forall xy\, (\phi_D(u, x) \wedge \psi_D(v, y))$.
3. $(\phi \vee \psi)^D \equiv \exists cuv \forall xy\, ((c = 0 \wedge \phi_D(u, x)) \vee (c \neq 0 \wedge \psi_D(v, y)))$.
4. $(\phi \supset \psi)^D \equiv \exists xv \forall uy\, (\phi_D(u, xuy) \supset \psi_D(vu, y))$.
5. $(\forall z\, \phi)^D \equiv \exists u \forall zx\, \phi_D(uz, x, z)$.
6. $(\exists z\, \phi)^D \equiv \exists zu \forall x\, \phi_D(u, x, z)$.

2.5. Theorem. [Gödel] *If ϕ is provable in HA, then there exist terms $\boldsymbol{u}$ such that $\boldsymbol{u}$ contains no negative variables in ϕ^D and $\phi_D(\boldsymbol{u}, x)$ is provable in T.*

The terms $\boldsymbol{u}$ will be called witnesses. The theorem is often formulated so that the terms $\boldsymbol{u}$ are closed terms and $\phi_D(\boldsymbol{u}z, x, z)$ holds. Assuming that there is a closed term for each type we can easily convert our witnesses into this form.

2.6. The analysis of Gödel's D-translation. The most important case in Gödel's D-translation is $(\phi \supset \psi)^D$. It quickly yields a very complicated formula due to the presence of u in xuy within the antecedent of $\phi_D(u, xuy) \supset \psi_D(vu, y)$. We analyze why u is necessary in xuy or why x needs to depend on u, and find that there are three reasons.

1. To obtain the logical equivalence of $\phi \supset (\psi \supset \tau)$ and $(\phi \wedge \psi) \supset \tau$.

2. To make x the choice functions of u in $\exists u \forall x\, (\phi(u, x) \supset \psi)$.

3. To obtain the validity of $\phi \supset \phi \wedge \phi$, i.e., contraction.

This leads us to the idea that if we move to a contraction-free logic, and supply necessary dependency for conjunction and existential quantifier, we can dispense with u in xuy.

Guided by the above idea we will modify the D-translation to the following which we distinguish from Gödel's by the superscript and subscript L.

2'. $(\phi \wedge \psi)^L \equiv \exists uv \forall xy\, (\phi_L(u, xv) \wedge \psi_L(v, yu))$.

4'. $(\phi \supset \psi)^L \equiv \exists xv \forall uy\, (\phi_L(u, xy) \supset \psi_L(vu, y))$.

6'. $(\exists z\, \phi)^L \equiv \exists zu \forall x\, \phi_L(u, xu, z)$.

The key difference from Gödel's original D-translation is that we erase the dependency of positive x on negative u in implication and raise instead the types of negative variables in conjunction and existential quantification.

2.7. De Paiva's dialectica category $\mathbb{GC}$. Inspired by Gödel's D-interpretation de Paiva defined two categories $\mathbb{DC}$ and $\mathbb{GC}$ [de Paiva 1991]. It turned out that the propositional part of our D-translation uses the same construction as $\mathbb{GC}$. We note that $\mathbb{GC}$ is very closely related to Chu's general construction of $*$-autonomous category [de Paiva 2006, Devarajan et al. 1999].

An object of $\mathbb{GC}$ is a relation, i.e. an equivalence class of monics $A \overset{\alpha}{\rightarrowtail} U \times X$, in a base category $\mathbb{C}$ which is required to be finitely complete.

A morphism from $A \overset{\alpha}{\rightarrowtail} U \times X$ to $B \overset{\beta}{\rightarrowtail} V \times Y$ in $\mathbb{GC}$ is a pair (f, F) of morphisms $f : U \to V$ and $F : Y \to X$ in $\mathbb{C}$ such that there exists a morphism k which makes the diagram

$$\begin{tikzcd}
& A' \arrow[r] \arrow[d, tail] \arrow[dl, "k"'] \arrow[dr, phantom, "\lrcorner", very near start] & A \arrow[d, tail, "\alpha"] \\
B' \arrow[r, tail] \arrow[dr, phantom, "\lrcorner", very near start] & U \times Y \arrow[r, "\mathtt{id}_U \times F"'] & U \times X \\
B \arrow[u] \arrow[r, tail, "\beta"'] & V \times Y \arrow[u, "f \times \mathtt{id}_Y"'] &
\end{tikzcd}$$

commute in $\mathbb{C}$, where the upper right and the lower left squares are pullback squares.

The category $\mathbb{GC}$ can be formulated more abstractly in terms of a pre-ordered fibration $p : \mathbb{P} \to \mathbb{C}$ as Hyland does with respect to $\mathbb{DC}$ [Hyland 2002]. We regard

$$\mathbb{P}(I) = (\{\alpha \in \mathbb{P} \mid p(\alpha) = I\}, \vdash_I)$$

as a pre-ordered collection of the predicates on I for each object I of $\mathbb{C}$. An Object of $\mathbb{GC}$ is then (α, U, X) with $\alpha \in \mathbb{P}(U \times X)$ and a morphism of $\mathbb{GC}$ from (α, U, X) to (β, V, Y) is (f, F) with $f : U \to V$ and $F : Y \to X$ in $\mathbb{C}$ such that

$$\alpha_{\mathtt{id}_U \times F} \vdash_{U \times Y} \beta_{f \times \mathtt{id}_Y}$$

holds for the Cartesian liftings $\alpha_{\mathtt{id}_U \times F} \to \alpha$ of α along $\mathtt{id}_U \times F$ and $\beta_{f \times \mathtt{id}_Y} \to \beta$ of β along $f \times \mathtt{id}_Y$.

We will however focus entirely on the concrete case with $\mathbb{C} = \mathbb{SET}$ since this case is most perspicuous and sufficient to motivate our work. An object is then simply a subset of $U \times X$ for a pair of sets U and X. Let us write $\alpha(u, x)$ for $(u, x) \in \alpha$. A morphism from $\alpha \subseteq U \times X$ to $\beta \subseteq V \times Y$ is then a pair (f, F) of functions $f : X \to Y$ and $F : V \to U$ such that

$$\{(u, y) \in U \times Y \mid \alpha(u, Fy)\} \subseteq \{(u, y) \in U \times Y \mid \beta(fu, y)\}.$$

holds.

The tensor product $\alpha \otimes \beta$ of $\alpha \subseteq U \times X$ and $\beta \subseteq V \times Y$ is defined as

$$\alpha \otimes \beta = \{(u, v, x, y) \in (U \times V) \times \left(X^V \times Y^U\right) \mid \alpha(u, xv) \wedge \beta(v, yu)\}.$$

and the internal hom object $\alpha \multimap \beta$ is

$$\alpha \multimap \beta = \{(x, v, u, y) \in \left(X^Y \times V^U\right) \times (U \times Y) \mid \alpha(u, xy) \supset \beta(vu, y)\}.$$

Furthermore the linear negation $\alpha^\perp$ of $\alpha \subseteq U \times X$ can be defined just as

$$\alpha^\perp = \{(x, u) \in X \times U \mid \neg\alpha(u, x)\}$$

and then $\alpha \parr \beta = \left(\alpha^\perp \otimes \beta^\perp\right)^\perp$ is

$$\alpha \parr \beta = \{(u, v, x, y) \in \left(U^Y \times V^X\right) \times (X \times Y) \mid \alpha(uy, x) \vee \beta(vx, y)\}.$$

The additive product $\alpha \& \beta$ and coproduct $\alpha \oplus \beta$ can also be defined.

For the exponential $!\,\alpha$ we have two choices. One is to define

$$!\,\alpha = \{(u, x) \in U \times (X^*)^U \mid \forall i \leq \ell(xu).\ \alpha(u, \mathtt{proj}_i(xu))\}$$

where the set X^* is the set of finite sequences of elements of X and $\ell(s)$ is the length of $s \in X^*$, and the other is

$$!\,\alpha = \{(u, x) \in U \times X^U \mid \alpha(u, xu)\}.$$

De Paiva uses the former definition while the latter is also indicated. As we will see later the former corresponds to the Diller-Nahm variant of D-translation and the latter to Gödel's original.

3. The Dialectica interpretation of first-order classical affine logic

3.1. The first-order classical affine logic. In the Gentzen-style sequent calculus of classical logic we have the structural rules

$$\frac{\vdash \Gamma, \phi, \psi, \Delta}{\vdash \Gamma, \psi, \phi, \Delta} \qquad \frac{\vdash \Gamma, \phi, \phi}{\vdash \Gamma, \phi} \qquad \frac{\vdash \Gamma}{\vdash \Gamma, \phi}$$

called exchange, contraction and weakening, respectively, where Γ and Δ are lists of formulas. By restricting the use of some of them we obtain a family of logical systems, which are generally called substructural logics.

Linear logic is a kind of substructural logic introduced by Girard, in which we have the propositional operators $?$ and $!$, and the use of contraction and weakening is restricted to formulas of the form $?\,\phi$.

For the computational aspect of logic the contraction rule is critical while weakening is often turned out not so much. For our work weakening is also admissible. We hence accept the weakening rule from the beginning and consider the first-order classical affine logic where the weakening rule

$$\frac{\vdash \Gamma}{\vdash \Gamma, \,?\,\phi}$$

of linear logic is replaced by the weakening for any formula ϕ. As a result of this we no longer need to distinguish multiplicative and additive constants.

We use the one-sided sequent calculus where cut has the form

$$\frac{\vdash \Gamma, \phi \qquad \vdash \Delta, \phi^{\perp}}{\vdash \Gamma, \Delta}$$

while the linear negation of ϕ is defined by the DeMorgan duality. For the standard formulation of linear logic we refer the reader to Girard's original paper [Girard 1987]. The following should however be noted.

1. We assume a collection of the pair of atomic n-ary predicate symbols $\mathtt{P}$ and $\bar{\mathtt{P}}$ such that $(\mathtt{P}(t_1, t_2, \ldots, t_n))^{\perp} \equiv \bar{\mathtt{P}}(t_1, t_2, \ldots, t_n)$ and $(\bar{\mathtt{P}}(t_1, t_2, \ldots, t_n))^{\perp} \equiv \mathtt{P}(t_1, t_2, \ldots, t_n)$. The atomic formulas with $\mathtt{P}$ are called positive while those with $\bar{\mathtt{P}}$ are negative.

2. We do not consider the propositional constants, not for any difficulty but for simplicity.

3. We assume the use of exchange rule freely and implicitly.

3.2. The equational system S of lambda calculus with the conditional. We consider the equational system of simply typed lambda calculus with the conditional constructor. The system will be called S for the time being. The types are constructed inductively from the single base type $\mathtt{0}$ by the operation $(\sigma, \tau) \mapsto \sigma \to \tau$. We assume that there are infinitely many variables and regard the pair (x, σ) of a variable x and a type σ as a variable of type σ. We write x^{σ} for (x, σ). The set of terms is defined as the smallest set satisfying

- any variable of σ is a term of σ,
- if $\mathtt{f}$ is an n-ary function symbol of the first-order affine logic then $\mathtt{f}$ is a term of $\underbrace{\mathtt{0} \to \cdots \to \mathtt{0}}_{n} \to \mathtt{0}$,
- 0 and 1 are terms of $\mathtt{0}$,
- $\mathtt{E}$ is a term of $\mathtt{0} \to \mathtt{0} \to \mathtt{0}$,
- if $\mathtt{P}$ is an n-ary positive predicate symbol of the first-order affine logic then $\mathtt{P_f}$ is a term of $\underbrace{\mathtt{0} \to \cdots \to \mathtt{0}}_{n} \to \mathtt{0}$,
- If s is a term of $\tau \to \sigma$ and t is a term of τ then st is a term of σ,
- $\mathtt{K}_{\sigma,\tau}$ is a term of $\sigma \to \tau \to \sigma$ for each pair of σ and τ,
- $\mathtt{S}_{\rho,\sigma,\tau}$ is a term of $(\rho \to \sigma \to \tau) \to (\rho \to \sigma) \to (\rho \to \tau)$ for each triple of ρ, σ and τ,
- $\mathtt{Cond}_\sigma$ is a term of $\mathtt{0} \to \sigma \to \sigma \to \sigma$ for each σ.

We will suppress the reference to types as much as possible by assuming that terms have always the matching types. From $\mathtt{S}$ and $\mathtt{K}$ we can define the lambda abstraction $\lambda x.\, t$ as usual. $\mathtt{E}$ and $\mathtt{P_f}$ are intended as the characteristic functions of the equality $=$ and the predicate $\mathtt{P}$, respectively.

The set of formulas is defined as the smallest set satisfying

- if $\mathtt{P}$ is a positive n-ary predicate symbol of the first-order affine logic and $t_1, t_2, \ldots, t_n$ are terms of type $\mathtt{0}$ then
$$\mathtt{P}(t_1, t_2, \ldots, t_n)$$
is a formula,
- if s and t are terms of type $\mathtt{0}$ then $s = t$ is a formula,
- if α and β are formulas then $\neg\alpha$, $\alpha \wedge \beta$, $\alpha \vee \beta$ and $\alpha \supset \beta$ are formulas.

Formulas with no $\neg$, $\wedge$, $\vee$, $\supset$ are called atomic. We write $s \neq t$ for $\neg s = t$.

We use the axioms and inference rules of classical propositional logic and equality. In addition we have the axioms

$$\vdash 0 \neq 1, \qquad \vdash \mathtt{K}(x, y)w = xw, \qquad \vdash \mathtt{S}(x, y, z)w = xz(yz)w,$$

$$\vdash x = 0 \supset \mathtt{Cond}(x, y_1, y_2)w = y_1 w, \qquad \vdash x \neq 0 \supset \mathtt{Cond}(x, y_1, y_2)w = y_2 w,$$

$$\vdash x = y \leftrightarrow \mathtt{E}(x, y) = 0, \qquad \vdash P(x_1, x_2, \ldots, x_n) \leftrightarrow \mathtt{P_f}(x_1, x_2, \ldots, x_n) = 0$$

with fresh variables w, and the inference rules

$$\frac{\vdash \phi}{\vdash \phi[t/x]} \qquad \frac{\vdash \phi[s/y] \quad \vdash sx = tx}{\vdash \phi[t/y]}$$

for any terms s and t of the matching type. The first inference rule is the rule of substitution and the second is the rule of weakly extensional equality between terms of a higher type in the style of Spector. For a higher-type terms s and t we use the expression $s = t$ only as the abbreviation to mean that $sx = tx$ is provable with an arbitrary x.

Note that we have the characteristic function $\lceil \alpha \rceil (u, x, z)$ available for any formula $\alpha(u, x, z)$ since we can define all the classical propositional functions from the conditional constructor. The negation is

$$\lambda x.\, \mathtt{Cond}(x, 1, 0)$$

and the conjunction is

$$\lambda xy.\, \mathtt{Cond}(x, (\mathtt{Cond}(y, 0, 1)), 1).$$

3.3. D-translation of first-order affine logic. For a formula ϕ of the first-order classical affine logic we inductively define its D-translation ϕ^L together with the quantifier free formula ϕ_L such that $\phi^L \equiv \exists u \forall x\, \phi_L(u, x)$. For formulas ϕ and ψ appearing in the inductive clauses we assume the renaming of bound variables so that

$$\phi^L \equiv \exists u \forall x\, \phi_L(u, x), \qquad \psi^L \equiv \exists v \forall y\, \psi_L(v, y)$$

where there is no overlapping of variables among u, v, x, y and free variables in ϕ^L or ψ^L.

3.4. Definition.

- $P^L \equiv P_L \equiv P$ *for a positive atomic formula* $P \equiv \mathtt{P}(t_1, \ldots, t_n)$.
- $Q^L \equiv Q_L \equiv \neg \mathtt{P}(t_1, \ldots, t_n)$ *for a negative atomic formula* $Q \equiv \bar{\mathtt{P}}(t_1, \ldots, t_n)$.
- $(\phi \otimes \psi)^L \equiv \exists uv \forall xy\, (\phi \otimes \psi)_L \equiv \exists uv \forall xy\, [\phi_L(u, xv) \wedge \psi_L(v, yu)]$.
- $(\phi \invamp \psi)^L \equiv \exists uv \forall xy\, (\phi \invamp \psi)_L \equiv \exists uv \forall xy\, [\phi_L(uy, x) \vee \psi_L(vx, y)]$.
- $(\phi \& \psi)^L \equiv \exists uv \forall cxy\, (\phi \& \psi)_L \equiv \exists uv \forall cxy\, [(c = 0 \supset \phi_L(u, x)) \wedge (c \neq 0 \supset \psi_L(v, y))]$.
- $(\phi \oplus \psi)^L \equiv \exists cuv \forall xy\, (\phi \oplus \psi)_L \equiv \exists cuv \forall xy\, [(c = 0 \wedge \phi_L(u, x)) \vee (c \neq 0 \wedge \psi_L(v, y))]$.
- $(\forall z\, \phi(z))^L \equiv \exists u \forall xz\, (\forall z\, \phi(z))_L \equiv \exists u \forall xz\, \phi_L(uz, x, z)$.
- $(\exists z\, \phi(z))^L \equiv \exists uz \forall x\, (\exists z\, \phi(z))_L \equiv \exists uz \forall x\, \phi_L(u, xz, z)$.
- $(\,!\,\phi)^L \equiv \exists u \forall x\, (\,!\,\phi)_L \equiv \exists u \forall x\, \phi_L(u, xu)$.
- $(\,?\,\phi)^L \equiv \exists u \forall x\, (\,?\,\phi)_L \equiv \exists u \forall x\, \phi_L(ux, x)$.

3.5. Lemma. *Suppose* $\phi^L \equiv \exists u \forall x\, \phi_L(u, x)$. *Then*

$$\left(\phi^\perp\right)^L \equiv \exists x \forall u \left(\phi^\perp\right)_L(u, x) \equiv \exists x \forall u\, \neg\phi_L(u, x).$$

Proof. The proof is by induction on the construction of affine logic formulas. The atomic case is obvious. The inductive cases are shown as follows.

$$\begin{aligned}
\left[(\phi \otimes \psi)^\perp\right]^L &\equiv \left(\phi^\perp \parr \psi^\perp\right)^L \\
&\equiv \exists xy \forall uv \left[\left(\phi^\perp\right)_L(u, xv) \vee \left(\psi^\perp\right)_L(v, yu)\right] \\
&\equiv \exists xy \forall uv \left[\neg\phi_L(u, xv) \vee \neg\psi_L(v, yu)\right] \\
&\equiv \exists xy \forall uv\, \neg \left[\phi_L(u, xv) \wedge \psi_L(v, yu)\right].
\end{aligned}$$

$$\begin{aligned}
\left[(\phi \& \psi)^\perp\right]^L &\equiv \left(\phi^\perp \oplus \psi^\perp\right)^L \\
&\equiv \exists cxy \forall uv \left[\left(c = 0 \wedge \left(\phi^\perp\right)_L(u, x)\right) \vee \left(c \neq 0 \wedge \left(\psi^\perp\right)_L(v, y)\right)\right] \\
&\equiv \exists cxy \forall uv \left[(c = 0 \wedge \neg\phi_L(u, x)) \vee (c \neq 0 \wedge \neg\psi_L(v, y))\right] \\
&\equiv \exists cxy \forall uv\, \neg \left[(c = 0 \supset \phi_L(u, x)) \wedge (c \neq 0 \supset \psi_L(v, y))\right].
\end{aligned}$$

$$\begin{aligned}
\left[(\,!\,\phi)^\perp\right]^L &\equiv \left(\,!\,\phi^\perp\right)^L \\
&\equiv \exists x \forall u \left(\phi^\perp\right)_L(ux, x) \\
&\equiv \exists x \forall u\, \neg\phi_L(ux, x).
\end{aligned}$$

The cases for $\phi \parr \psi, \phi \oplus \psi$ and $?\,\phi$ are shown similarly. ∎

3.6. The Dialectica interpretation of first-order classical affine logic. We will prove the theorem analogous to Theorem 2.5 for the first-order classical affine logic using ϕ^L and S instead of ϕ^D and T.

3.7. Theorem. *If ϕ is provable in first-order classical affine logic, then there exist terms $\boldsymbol{u}$ such that $\boldsymbol{u}$ contains no negative variables in ϕ^L and $\phi_L(\boldsymbol{u}, x)$ is provable in S.*

The proof is by induction on the construction of proofs. We suppress the variables z in $\alpha(u, x, z)$ as much as possible as before. First let us prepare some preliminary lemmas.

3.8. Lemma. *Let $\phi \equiv \phi_1 \parr \phi_2$ and $\psi \equiv \phi_2 \parr \phi_1$. Then $\phi_L(u, x) \leftrightarrow \psi_L(u, x)$ holds.*

Proof. Let $(\phi_i)^L \equiv \exists u_i \forall x_i (\phi_i)_L(u_i, x_i)$. Then $\phi_L \equiv (\phi_1)_L(u_1 x_2, x_1) \vee (\phi_2)_L(u_2 x_1, x_2)$ and $\psi_L \equiv (\phi_2)_L(u_2 x_1, x_2) \vee (\phi_1)_L(u_1 x_2, x_1)$. ∎

3.9. Lemma. *Let* $\phi \equiv (\phi_1 ⅋ \phi_2) ⅋ \phi_3$ *and* $\psi \equiv \phi_1 ⅋ (\phi_2 ⅋ \phi_3)$. *There exist closed terms* t_i *and* s_i *such that*

$$\begin{cases} \phi_L(u_1, u_2, \dots, u_n, x) \leftrightarrow \psi_L(t_1 u_1, t_2 u_2, \dots, t_n u_n, x) \\ \phi_L(s_1 u_1, s_2 u_2, \dots, s_n u_n, x) \leftrightarrow \psi_L(u_1, u_2, \dots, u_n, x) \end{cases}$$

hold where $u_1, u_2, \dots, u_n$ *is the list of all positive variables.*

Proof. For each positive u_{ik} of $(\phi_i)_L$ we have $(\phi_i)_L(u_{ik}y, x_i)$ in ϕ_L and $(\phi_i)_L(u_{ik}z, x_i)$ in ψ_L where y and z are permutations of negative x_j's with $i \neq j$. Let $t_{ik} \equiv \lambda uz.\, uy$ and $s_{ik} \equiv \lambda uy.\, uz$. ∎

3.10. Lemma. *Let* $\phi^L \equiv \exists u \forall x\, \phi_L(u, x)$ *and* $\psi^L \equiv \exists v \forall y\, \psi_L(v, y)$. *There exist closed terms* t_i *and* s_i *such that*

$$\begin{cases} (\,?\,\phi ⅋ \psi)_L(u_1, \dots, u_n, v, x, y) \leftrightarrow \phi_L(t_1 u_1 z_1, \dots, t_n u_n z_n, x) \vee \psi_L(vx, y) \\ (\,?\,\phi ⅋ \psi)_L(s_1 u_1, \dots, s_n u_n, v, x, y) \leftrightarrow \phi_L(u_1 z_1, \dots, u_n z_n, x) \vee \psi_L(vx, y) \end{cases}$$

hold where z_i *is a permutation of* x, y *and* $u_1, u_2, \dots, u_n$ *is the list* u.

Proof. Recall $(\,?\,\phi ⅋ \psi)_L \equiv \phi_L(u_1 yx, \dots, u_n yx, x) \vee \psi_L(vx, y)$. Let $t_i \equiv \lambda u z_i.\, uyx$ and $s_i \equiv \lambda yx.\, u z_i$. ∎

Lemma 3.9 takes care of the parenthesizing with respect to ⅋ and lets us handle the list Γ as a single formula. Lemma 3.8 together with Lemma 3.9 justifies the implicit use of exchange. Lemma 3.10 is used for the promotion rule. With this preparation in mind let us begin the proof of Theorem 3.7.

Axioms.

$$\vdash \phi, \phi^\perp$$

Let $\phi^L \equiv \exists u \forall x\, \phi_L(u, x)$. By Lemma 3.5 we have $\left(\phi^\perp\right)^L \equiv \exists x \forall u\, \neg\phi_L(u, x)$. Renaming variables we need to find $\boldsymbol{u}$ and $\boldsymbol{x}$ such that

$$\phi_L(\boldsymbol{u}v, y) \vee \neg\phi_L(v, \boldsymbol{x}y)$$

holds. The identities $\lambda v.\, v$ and $\lambda y.\, y$ suffice for them.

$\otimes$-rule.

$$\frac{\vdash \Gamma, \phi \qquad \vdash \Delta, \psi}{\vdash \Gamma, \Delta, \phi \otimes \psi}$$

Let

$$\Gamma^L \equiv \exists w_1 \forall z_1\, \Gamma_L(w_1, z_1), \quad \phi^L \equiv \exists u \forall x\, \phi_L(u, x),$$
$$\Delta^L \equiv \exists w_2 \forall z_2\, \Delta_L(w_2, z_2), \quad \psi^L \equiv \exists v \forall y\, \psi_L(v, y).$$

By the inductive hypothesis we have already found the witness terms $\boldsymbol{u}, \boldsymbol{v}, \boldsymbol{w_1}$ and $\boldsymbol{w_2}$ such that we can prove

$$\Gamma_L(\boldsymbol{w_1}x, z_1) \vee \phi_L(\boldsymbol{u}z_1, x) \qquad \text{and} \qquad \Delta_L(\boldsymbol{w_2}y, z_2) \vee \psi_L(\boldsymbol{v}z_2, y),$$

and we need to find witnesses $\widetilde{\boldsymbol{u}}, \widetilde{\boldsymbol{v}}, \widetilde{\boldsymbol{w}}_1$ and $\widetilde{\boldsymbol{w}}_2$ such that

$$\Gamma_L(\widetilde{\boldsymbol{w_1}}xyz_2, z_1) \vee \Delta_L(\widetilde{\boldsymbol{w_2}}xyz_1, z_2) \vee [\phi_L(\widetilde{\boldsymbol{u}}z_1z_2, x(\widetilde{\boldsymbol{v}}z_1z_2)) \wedge \psi_L(\widetilde{\boldsymbol{v}}z_1z_2, y(\widetilde{\boldsymbol{u}}z_1z_2))]$$

holds. We may choose them so that we have

$$\widetilde{\boldsymbol{u}}z_1z_2 = \boldsymbol{u}z_1, \quad \widetilde{\boldsymbol{v}}z_1z_2 = \boldsymbol{v}z_2$$

and

$$\widetilde{\boldsymbol{w_1}}xyz_2 = \boldsymbol{w_1}(x(\widetilde{\boldsymbol{v}}z_1z_2)) = \boldsymbol{w_1}(x(\boldsymbol{v}z_2)), \quad \widetilde{\boldsymbol{w_2}}xyz_1 = \boldsymbol{w_2}(y(\widetilde{\boldsymbol{u}}z_1z_2)) = \boldsymbol{w_2}(y(\boldsymbol{u}z_1)).$$

where x and y have the types assigned in the lower sequent. From the inductive hypothesis we can prove

$$\Gamma_L(\boldsymbol{w_1}(x(\boldsymbol{v}z_2)), z_1) \vee \phi_L(\boldsymbol{u}z_1, x(\boldsymbol{v}z_2)) \quad \text{and} \quad \Delta_L(\boldsymbol{w_2}(y(\boldsymbol{u}z_1)), z_2) \vee \psi_L(\boldsymbol{v}z_2, y(\boldsymbol{u}z_1))$$

by substitution. The conclusion then follows by the propositional calculus.

⅋-rule.

$$\frac{\vdash \Gamma, \phi, \psi}{\vdash \Gamma, \phi ⅋ \psi}$$

The case reduces to Lemma 3.9.

&-rule.

$$\frac{\vdash \Gamma, \phi \qquad \vdash \Gamma, \psi}{\vdash \Gamma, \phi \& \psi}$$

For two copies of Γ in the upper sequents we let

$$\Gamma^L \equiv \exists w_1 \forall z_1\, \Gamma_L(w_1, z_1), \qquad \Gamma^L \equiv \exists w_1 \forall z_1\, \Gamma_L(w_2, z_2)$$

and for Γ in the lower sequent

$$\Gamma^L \equiv \exists w \forall z\, \Gamma_L(w, z).$$

Let $\phi^L \equiv \exists u \forall x\, \phi_L(u, x)$ and $\psi^L \equiv \exists v \forall y\, \psi_L(v, y)$. By inductive hypothesis there exist $\boldsymbol{u}, \boldsymbol{v}, \boldsymbol{w_1}$ and $\boldsymbol{w_2}$ such that

$$\Gamma_L(\boldsymbol{w_1}x, z_1) \vee \phi_L(\boldsymbol{u}z_1, x), \qquad \Gamma_L(\boldsymbol{w_2}y, z_2) \vee \psi_L(\boldsymbol{v}z_2, y)$$

hold. We need to find $\widetilde{\boldsymbol{u}}, \widetilde{\boldsymbol{v}}$ and $\widetilde{\boldsymbol{w}}$ such that

$$\Gamma_L(\widetilde{\boldsymbol{w}}cxy, z) \vee [(c = 0 \supset \phi_L(\widetilde{\boldsymbol{u}}z, x)) \wedge (c \neq 0 \supset \psi_L(\widetilde{\boldsymbol{v}}z, y))]$$

holds. We choose $\widetilde{\boldsymbol{u}}$ and $\widetilde{\boldsymbol{v}}$ so that we have $\widetilde{\boldsymbol{u}}z = \boldsymbol{u}z$ and $\widetilde{\boldsymbol{v}}z = \boldsymbol{v}z$, and choose $\widetilde{\boldsymbol{w}}$ according to c so that

$$\widetilde{\boldsymbol{w}}cxy = \begin{cases} \boldsymbol{w_1}x & \text{if } c = 0 \\ \boldsymbol{w_2}y & \text{otherwise} \end{cases}$$

holds. Since we have the conditional we can in fact extract $\widetilde{\boldsymbol{w}}$ from the equation

$$\widetilde{\boldsymbol{w}}cxy = \mathtt{Cond}(c, \boldsymbol{w_1}x, \boldsymbol{w_2}y).$$

The conclusion then follows from the inductive hypothesis by the propositional calculus.

$\oplus$-rule.

$$\frac{\vdash \Gamma, \phi_i}{\vdash \Gamma, \phi_1 \oplus \phi_2} \quad (i = 1, 2)$$

For ϕ_i in the upper sequent we let $(\phi_i)^L \equiv \exists u \forall x\, (\phi_i)_L(u, x)$ and for ϕ_1 and ϕ_2 in the lower sequent

$$(\phi_1)^L \equiv \exists u_1 \forall x_1\, (\phi_1)_L(u_1, x_1), \qquad (\phi_2)^L \equiv \exists u_2 \forall x_2\, (\phi_2)_L(u_2, x_2).$$

Let $\Gamma^L \equiv \exists v \forall y\, \Gamma_L(v, y)$. By the inductive hypothesis we have $\boldsymbol{u}$ and $\boldsymbol{v}$ such that

$$\Gamma_L(\boldsymbol{v}x, y) \vee (\phi_i)_L(\boldsymbol{u}y, x)$$

holds. We need to find $\widetilde{\boldsymbol{c}}, \widetilde{\boldsymbol{u_1}}, \widetilde{\boldsymbol{u_2}}$ and $\widetilde{\boldsymbol{v}}$ such that

$$\Gamma_L(\widetilde{\boldsymbol{v}}x_1x_2, y) \vee [(\widetilde{\boldsymbol{c}}y = 0 \wedge (\phi_1)_L(\widetilde{\boldsymbol{u_1}}y, x_1)) \vee (\widetilde{\boldsymbol{c}}y \neq 0 \wedge (\phi_2)_L(\widetilde{\boldsymbol{u_2}}y, x_2))]$$

holds. If $i = 1$ we choose $\widetilde{\boldsymbol{c}}, \widetilde{\boldsymbol{u_1}}$ and $\widetilde{\boldsymbol{v}}$ so that we have

$$\widetilde{\boldsymbol{c}}y = 0, \quad \widetilde{\boldsymbol{u_1}}y = \widetilde{\boldsymbol{u}}y, \quad \widetilde{\boldsymbol{v}}x_1x_2 = \boldsymbol{v}x_1$$

while $\widetilde{\boldsymbol{u_2}}$ can be any term of the matching type. If $i = 2$ we choose $\widetilde{\boldsymbol{c}}, \widetilde{\boldsymbol{u_2}}$ and $\widetilde{\boldsymbol{v}}$ so that we have

$$\widetilde{\boldsymbol{c}}y = 1, \quad \widetilde{\boldsymbol{u_2}}y = \widetilde{\boldsymbol{u}}y, \quad \widetilde{\boldsymbol{v}}x_1x_2 = \boldsymbol{v}x_2$$

while $\widetilde{\boldsymbol{u_1}}$ can be any term of the matching type. The conclusion then follows from the inductive hypothesis by the propositional calculus.

$\forall$-rule.

$$\frac{\vdash \Gamma, \phi}{\vdash \Gamma, \forall y\, \phi}$$

Let $\Gamma^L \equiv \exists w \forall z\, \Gamma_L(w, z)$ and $\phi^L \equiv \exists u \forall x\, \phi_L(u, x, y)$. By the inductive hypothesis we have $\boldsymbol{u}$ and $\boldsymbol{w}$ such that

$$\Gamma_L(\boldsymbol{w}x, z) \vee \phi_L(\boldsymbol{u}z, x, y)$$

holds, and we need to find $\widetilde{\boldsymbol{u}}$ and $\widetilde{\boldsymbol{w}}$ such that

$$\Gamma_L(\widetilde{\boldsymbol{w}}xy, z) \vee \phi_L(\widetilde{\boldsymbol{u}}yz, x, y)$$

holds. We choose them in such a way that we have

$$\widetilde{\boldsymbol{u}}yz = \boldsymbol{u}z, \qquad \widetilde{\boldsymbol{w}}xy = \boldsymbol{w}x$$

where $\widetilde{\boldsymbol{u}}$ and $\widetilde{\boldsymbol{w}}$ may contain y as a free variable. The conclusion then immediately follows from the hypothesis.

The eigenvariable condition is crucial since if Γ contained a free y we would have to rename y in $(\forall y\, \phi)^L$ resulting in

$$\Gamma_L(\widetilde{\boldsymbol{w}}xy', z, y) \vee \phi_L(\widetilde{\boldsymbol{u}}y'z, x, y')$$

which is no longer derivable from the hypothesis.

∃-rule.

$$\frac{\vdash \Gamma, \phi[t/v]}{\vdash \Gamma, \exists v\, \phi}$$

Let $\Gamma^L \equiv \exists w \forall z\, \Gamma_L(w, z)$ and $(\phi[t/v])^L \equiv \exists u \forall x\, (\phi[t/v])_L(u, x)$ where u and x are different from v and any free variable in t. We hence have

$$\exists u \forall x\, (\phi[t/v])_L(u, x) \equiv \exists u \forall x\, (\phi_L(u, x, v)[t/v]) \equiv \exists u \forall x\, \phi_L(u, x, t).$$

By the inductive hypothesis there exist $\boldsymbol{w}$ and $\boldsymbol{u}$ such that

$$\Gamma_L(\boldsymbol{w}x, z) \vee \phi_L(\boldsymbol{u}z, x, t)$$

holds. We need to find $\widetilde{\boldsymbol{w}}, \widetilde{\boldsymbol{u}}$ and $\widetilde{\boldsymbol{v}}$ such that

$$\Gamma_L(\widetilde{\boldsymbol{w}}x, z) \vee \phi_L(\widetilde{\boldsymbol{u}}z, x(\widetilde{\boldsymbol{v}}z), \widetilde{\boldsymbol{v}}z)$$

holds. It suffices to have

$$\widetilde{\boldsymbol{u}}z = \boldsymbol{u}z, \quad \widetilde{\boldsymbol{v}}z = t, \quad \widetilde{\boldsymbol{w}}x = \boldsymbol{w}(x(\widetilde{\boldsymbol{v}}z)) = \boldsymbol{w}(xt)$$

where x has the type in the lower sequent. Since t contains neither x nor z we can construct $\widetilde{\boldsymbol{v}}$ without using x, and $\widetilde{\boldsymbol{w}}$ without using z. The conclusion follows from the hypothesis by substitution.

Cut.

$$\frac{\vdash \Gamma, \phi \qquad \vdash \Delta, \phi^\perp}{\vdash \Gamma, \Delta}$$

Let

$$\Gamma^L \equiv \exists v \forall y\, \Gamma_L(v, y), \qquad \Delta^L \equiv \exists w \forall z\, \Delta_L(w, z),$$

$$\phi^L \equiv \exists u \forall x\, \phi_L(u, x).$$

Then $\left(\phi^\perp\right)^L \equiv \exists x \forall u\, \neg\phi_L(u, x)$ by Lemma 3.5. By the inductive hypothesis we have $\boldsymbol{v}, \boldsymbol{w}, \boldsymbol{u}$ and $\boldsymbol{x}$ such that

$$\Gamma_L(\boldsymbol{v}x, y) \vee \phi_L(\boldsymbol{u}y, x), \qquad \Delta_L(\boldsymbol{w}u, z) \vee \neg\phi_L(u, \boldsymbol{x}z)$$

both hold. We need to find $\widetilde{\boldsymbol{v}}$ and $\widetilde{\boldsymbol{w}}$ such that

$$\Gamma_L(\widetilde{\boldsymbol{v}}z, y) \vee \Delta_L(\widetilde{\boldsymbol{w}}y, z)$$

hold. Consider the substitution instances

$$\Gamma_L(\boldsymbol{v}(\boldsymbol{x}z), y) \vee \phi_L(\boldsymbol{u}y, \boldsymbol{x}z), \qquad \Delta_L(\boldsymbol{w}(\boldsymbol{u}y), z) \vee \neg\phi_L(\boldsymbol{u}y, \boldsymbol{x}z)$$

from the inductive hypothesis. It follows from them by the propositional calculus that

$$\Gamma_L(\boldsymbol{v}(\boldsymbol{x}z), y) \vee \Delta_L(\boldsymbol{w}(\boldsymbol{u}y), z).$$

We hence choose $\widetilde{\boldsymbol{v}}$ and $\widetilde{\boldsymbol{w}}$ so that we have

$$\widetilde{\boldsymbol{v}}z = \boldsymbol{v}(\boldsymbol{x}z), \qquad \widetilde{\boldsymbol{w}}y = \boldsymbol{w}(\boldsymbol{u}y).$$

The conclusion then holds immediately.

Weakening.

$$\frac{\vdash \Gamma}{\vdash \Gamma, \phi}$$

Let $\Gamma^L \equiv \exists v \forall y\, \Gamma_L(v, y)$ and $\phi^L \equiv \exists u \forall x\, \phi_L(u, x)$. By the inductive hypothesis there exists $\boldsymbol{v}$ such that

$$\Gamma_L(\boldsymbol{v}, y)$$

holds. We need to find $\widetilde{\boldsymbol{u}}$ and $\widetilde{\boldsymbol{v}}$ such that

$$\Gamma_L(\widetilde{\boldsymbol{v}}x, y) \vee \phi_L(\widetilde{\boldsymbol{u}}y, x)$$

holds. For this it suffices to have $\widetilde{\boldsymbol{v}}x = \widetilde{\boldsymbol{v}}$. Any term of the matching type would do for $\widetilde{\boldsymbol{u}}$.

Dereliction.

$$\frac{\vdash \Gamma, \phi}{\vdash \Gamma, ?\,\phi}$$

Let $\Gamma^L \equiv \exists v \forall y\, \Gamma_L(v, y)$ and $\phi^L \equiv \exists u \forall x\, \phi_L(u, x)$. By the inductive hypothesis there exist $\boldsymbol{u}$ and $\boldsymbol{v}$ such that

$$\Gamma_L(\boldsymbol{v}x, y) \vee \phi_L(\boldsymbol{u}y, x)$$

holds. We need to find $\widetilde{\boldsymbol{u}}$ and $\widetilde{\boldsymbol{v}}$ such that

$$\Gamma_L(\widetilde{\boldsymbol{v}}x, y) \vee \phi_L(\widetilde{\boldsymbol{u}}xy, x)$$

holds. It suffices to have $\widetilde{\boldsymbol{v}}x = \boldsymbol{v}x$ and $\widetilde{\boldsymbol{u}}xy = \boldsymbol{u}x$.

Contraction.

$$\frac{\vdash \Gamma, ?\,\phi, ?\,\phi}{\vdash \Gamma, ?\,\phi}$$

Let $\Gamma^L \equiv \exists v \forall y\, \Gamma_L(v, y)$. For two copies of $?\,\phi$ in the upper sequent we let

$$(?\,\phi)_L \equiv \exists u_1 \forall x_1\, \phi_L(u_1 x_1, x_1), \qquad (?\,\phi)_L \equiv \exists u_2 \forall x_2\, \phi_L(u_2 x_2, x_2)$$

respectively. For $?\,\phi$ in the lower sequent we let

$$(?\,\phi)^L \equiv \exists u \forall x\, \phi^L(ux, x).$$

By the inductive hypothesis we have $\boldsymbol{u_1}, \boldsymbol{u_2}$ and $\boldsymbol{v}$ such that

$$\Gamma_L(\boldsymbol{v}x_1x_2, y) \vee \phi_L(\boldsymbol{u_1}x_1x_2y, x_1) \vee \phi_L(\boldsymbol{u_2}x_1x_2y, x_2)$$

holds. We need to find $\widetilde{\boldsymbol{u}}$ and $\widetilde{\boldsymbol{v}}$ such that

$$\Gamma_L(\widetilde{\boldsymbol{v}}x, y) \vee \phi_L(\widetilde{\boldsymbol{u}}xy, x)$$

holds. To find them we first force x_1 and x_2 in the inductive hypothesis to the identical x by substitution so that

$$\Gamma_L(\boldsymbol{v}xx, y) \vee \phi_L(\boldsymbol{u_1}xxy, x) \vee \phi_L(\boldsymbol{u_2}xxy, x).$$

We then choose $\widetilde{\boldsymbol{u}}$ and $\widetilde{\boldsymbol{v}}$ in such a way that

$$\widetilde{\boldsymbol{v}}x = \boldsymbol{v}xx, \quad \widetilde{\boldsymbol{u}}xy = \begin{cases} \boldsymbol{u_1}xxy & \text{if } \phi_L(\boldsymbol{u_1}xxy, x) \\ \boldsymbol{u_2}xxy & \text{otherwise.} \end{cases}$$

Since we have the characteristic function $\lceil \phi_L(\boldsymbol{u_1}xxy, x) \rceil(x, y)$ of $\phi_L(\boldsymbol{u_1}xxy, x)$ with

$$\phi_L(\boldsymbol{u_1}xxy, x) \leftrightarrow \lceil \phi_L(\boldsymbol{u_1}xxy, x) \rceil(x, y) = 0$$

we can in fact extract $\widetilde{\boldsymbol{u}}$ from the equation

$$\widetilde{\boldsymbol{u}}xy = (\texttt{Cond}(\lceil \phi_L(\boldsymbol{u_1}xxy, x) \rceil(x, y), \boldsymbol{u_1}, \boldsymbol{u_2}))\, xyy.$$

The conclusion then follows by the propositional logic, using proof by cases.

The above tactics would not work if $?\phi$ was just ϕ. The inductive hypothesis would then become

$$\Gamma_L(\boldsymbol{v}x_1x_2, y) \vee \phi_L(\boldsymbol{u_1}x_2y, x_1) \vee \phi_L(\boldsymbol{u_2}x_1y, x_2)$$

and the conclusion

$$\Gamma_L(\widetilde{\boldsymbol{v}}x, y) \vee \phi_L(\widetilde{\boldsymbol{u}}y, x).$$

We would then not be able to choose between $\boldsymbol{u_1}xy$ and $\boldsymbol{u_2}xy$ in $\widetilde{\boldsymbol{u}}$ since x would not be available to $\widetilde{\boldsymbol{u}}$.

Promotion.

$$\frac{\vdash\ ?\Gamma, \phi}{\vdash\ ?\Gamma,\ !\phi}$$

Let $?\Gamma$ be the list $?\psi_1, ?\psi_2, \ldots ?\psi_n$ and $(\psi_i)^L \equiv \exists v_i \forall y_i\, (\psi_i)_L(v_i, y_i)$ for $1 \leq i \leq n$. By Lemma 3.10 we may write

$$(\,?\Gamma)^L \equiv \exists v \forall y\, [(\psi_1)_L(v_1 y, y_1) \vee (\psi_2)_L(v_2 y, y_2) \vee \cdots \vee (\psi_n)_L(v_n y, y_n)]$$

with v and y being the concatenations of $v_1, v_2, \ldots, v_n$ and $y_1, y_2, \ldots, y_n$, respectively. We write $\psi(z, y)$ for the quantifier free formula

$$(\psi_1)_L(z_1, y_1) \vee (\psi_2)_L(z_2, y_2) \vee \cdots \vee (\psi_n)_L(z_n, y_n).$$

with z being the concatenation of $z_1, z_2, \ldots, z_n$ so that $(\,?\Gamma)^L \equiv \exists v \forall y\, \psi(vy, y)$.

Let $\phi^L \equiv \exists u \forall x\, \phi_L(u, x)$. By the inductive hypothesis we have $\boldsymbol{u}$ and $\boldsymbol{v}$ such that

$$\psi(\boldsymbol{v}xy, y) \vee \phi_L(\boldsymbol{u}y, x)$$

holds, and we need to find $\widetilde{\boldsymbol{u}}$ and $\widetilde{\boldsymbol{v}}$ such that

$$\psi(\widetilde{\boldsymbol{v}}xy, y) \vee \phi_L(\widetilde{\boldsymbol{u}}y, x(\widetilde{\boldsymbol{u}}y))$$

holds. We choose $\widetilde{\boldsymbol{u}}$ and $\widetilde{\boldsymbol{v}}$ so that we have

$$\widetilde{\boldsymbol{u}}y = \boldsymbol{u}y, \qquad \widetilde{\boldsymbol{v}}xy = \boldsymbol{v}(x(\widetilde{\boldsymbol{u}}y))y = \boldsymbol{v}(x(\boldsymbol{u}y))y$$

where x has the type in the lower sequent. The conclusion then follows from the inductive hypothesis by substitution.

The above tactics would not work if $?\,\Gamma$ was just Γ. The inductive hypothesis would then become

$$\Gamma_L(\boldsymbol{v}x, y) \vee \phi_L(\boldsymbol{u}y, x)$$

and the conclusion

$$\Gamma_L(\widetilde{\boldsymbol{v}}x, y) \vee \phi_L(\widetilde{\boldsymbol{u}}y, x(\widetilde{\boldsymbol{u}}y)).$$

We would then not be able to choose $\widetilde{\boldsymbol{v}}$ so that $\widetilde{\boldsymbol{v}}x = \boldsymbol{v}(x(\boldsymbol{u}y))$ since some of the variables in y would not be available to $\widetilde{\boldsymbol{v}}$.

3.11. The Diller-Nahm variant of the Dialectica interpretation. In Gödel's original Dialectica interpretation we need first evaluate the truth value of a formula within a witness term for the contraction, in the same way as we have just seen. This causes no problem since the only predicate considered is the equality between natural numbers, whose characteristic function is primitive recursive.

The Dialectica interpretation can be naturally extended to higher-order systems. We then need to consider the equality between higher-type functionals, whose characteristic function is, however, neither continuous nor provably recursive. This observation motivated Diller and Nahm to give a variant of the Dialectica interpretation, in which the use of characteristic function is no longer necessary [Diller and Nahm 1974].

The Diller-Nahm variant of the Dialectica interpretation is also important from the viewpoint of categorical logic, as pointed out by Hyland [Hyland 2002]. The categorical counterpart of the standard Dialectica interpretation forms a symmetric monoidal closed category, while the Diller-Nahm variant immediately yields a Cartesian closed category.

We can modify our D-translation in the style of Diller-Nahm simply by adopting the definition $!\,\alpha = \{(u,x) \in U \times (X^*)^U \mid \forall i \leq \ell(xu).\ \alpha(u, \mathtt{proj}_i(xu))\}$ and its dual form for $?\,\alpha$ in de Paiva's $\mathbb{GC}$ [de Paiva 1991]. We first need to extend types so that we have the type σ^* for each type σ. Its intended interpretation $[\![\sigma^*]\!]$ is the free commutative monoid X^* generated from $[\![\sigma]\!] = X$. We also need to extend terms to represent the monoid unit $e \in X^*$, the monoid multiplication $\cdot : X^* \times X^* \to X^*$, the unit map $\eta_X : X \to X^*$ and the counit map $\epsilon_{X^*} : X^{**} \to X^*$. For $a \in X$ and $s \in X^*$ we write $a \in s$ if $s = s'\eta_X(a)$ for some $s' \in X^*$.

For the sake of simplicity we only consider the case where the formula α has a single positive variable u and a single negative variable x.

We define the operations $(\_)^\circ$ and $(\_)^\bullet$ on the formula $\alpha(u, v)$ by

$$\alpha^\circ(u, x) \qquad \text{iff} \qquad \text{for some } w \in u, \quad \alpha(w, x)$$

where the type of u is raised from $\tau \to \sigma$ to $\tau \to \sigma^*$, and

$$\alpha^\bullet(u, x) \qquad \text{iff} \qquad \text{for all } z \in x, \quad \alpha(u, z)$$

where the type of x is raised as before.

The modified translation ϕ^{DN} is the same as ϕ^{L} except for

- $(\,?\,\phi)^{DN} \equiv \exists u \forall x\ (\phi_{DN})^{\circ}\,(ux, x)$,
- $(\,!\,\phi)^{DN} \equiv \exists u \forall x\ (\phi_{DN})^{\bullet}\,(u, xu)$.

This gives us the translation

$$(\,!\,\phi \multimap \psi)^{DN} \equiv \exists xv \forall uy\,[(\phi_{DN})^{\bullet}\,(u, xyu) \supset \psi_{DN}(vu, y)]$$

as Hyland defined for the intuitionistic implication $\phi \supset \psi$.

3.12. The first-order extension of $\mathbb{GC}$. De Paiva's $\mathbb{GC}$ gives us a categorical model of propositional classical linear logic. We show how our D-translation extends $\mathbb{GC}$ to the first-order setting in terms of an indexed category. Such an extension is already given for variants of $\mathbb{DC}$ [Hyland 2002, Streicher 2000]. We however think that it is worthwhile looking into our construction in some detail since it is concrete and simple. The construction is quite general, but we again focus on the the case where $\mathbb{C} = \mathbb{SET}$.

We first define the category $\mathbb{GC}(Z)$ for each object Z in $\mathbb{C}$. An object is a subset of $U \times X \times Z$ for a pair of sets U and X with the fixed Z.

$$\{(u, x, z) \in U \times X \times Z \mid \alpha(u, x, z)\}.$$

A morphism from $\alpha \subseteq U \times X \times Z$ to $\beta \subseteq V \times Y \times Z$ is a pair (f, F) of morphisms $f : U \times Z \to V$ and $F : Y \times Z \to X$ such that

$$\{(u, y, z) \in U \times Y \times Z \mid \alpha(u, Fyz, z)\} \subseteq \{(u, y, z) \in U \times Y \times Z \mid \beta(fuz, y, z)\}$$

holds. The composition of a morphism (f, F) from α to β and a morphism (g, G) from β to γ is then defined as $(g \circ \langle f, \mathtt{proj}_2\rangle, F \circ \langle G, \mathtt{proj}_2\rangle)$. It amounts to derive

$$\alpha(u, F(Gz'z)z, z) \supset \gamma(g(fuz)z, z', z)$$

from

$$\begin{cases} \alpha(u, Fyz, z) \supset \beta(fuz, y, z) \\ \beta(v, Gz'z, z) \supset \gamma(gvz, z', z) \end{cases}$$

by substituting fuz for v and $Gz'z$ for y.

The $*$-autonomous structures of $\mathbb{GC}$ can be naturally extended to $\mathbb{GC}(Z)$ by defining

$$\begin{aligned} \alpha^{\perp} &= \{(x, u, z) \in X \times U \times Z \mid \neg\alpha(u, x, z)\}, \\ \alpha \otimes \beta &= \{(u, v, x, y, z) \mid (U \times V) \times (X^{V} \times Y^{U}) \times Z \mid \alpha(u, xv, z) \wedge \beta(v, yu, z)\} \end{aligned}$$

for $\alpha \subseteq U \times X \times Z$ and $\beta \subseteq V \times Y \times Z$. The monoidal closedness is established by the bijection between $(f, \langle F_1, F_2\rangle)$ and $(\langle \tilde{f}_1, \tilde{f}_2\rangle, \tilde{F})$ in

$$\begin{aligned} &\alpha(u, (F_1 z'z)v, z) \wedge \beta(v, (F_2 z'z)u, z) \supset \gamma(fuvz, z', z), \\ &\alpha(u, \tilde{F}vz'z, z) \supset \beta(v, (\tilde{f}_1 uz)z', z) \supset \gamma((\tilde{f}_2 uz)v, z', z). \end{aligned}$$

Furthermore we have a functor f^{-1} from $\mathbb{GC}(Z)$ to $\mathbb{GC}(Z')$ for any $f : Z' \to Z$ in $\mathbb{C}$ defined by

$$\begin{cases} f^{-1}(\alpha) = \{(u, x, z') \in U \times X \times Z' \mid \alpha(u, x, fz')\}, \\ f^{-1}(g, G) = (g \circ (\mathtt{id}_U \times f), G \circ (\mathtt{id}_Y \times f)) \end{cases}$$

for $\alpha \subseteq U \times X \times Z$ and $g : U \times Z \to V,\ G : Y \times Z \to X$. It just changes $\alpha(u, Gyz, z) \supset \beta(guz, y, z)$ into

$$\alpha(u, Gy(fz'), fz') \supset \beta(gu(fz'), y, fz').$$

We hence established that the pair of maps

$$\begin{cases} Z \mapsto \mathbb{GC}(Z), \\ f \mapsto f^{-1} \end{cases}$$

is a contravariant functor from $\mathbb{C} = \mathbb{SET}$ to $\mathbb{CAT}$, i.e. an indexed category.

The quantifiers $\forall$ and $\exists$ are functors from $\mathbb{GC}(Z \times Z')$ to $\mathbb{GC}(Z')$ defined according to our D-translation. The functor $\forall$ is given by

$$\begin{cases} \forall(\alpha) = \{(u, x, z, z') \in U^Z \times X \times Z \times Z' \mid \alpha(uz, x, z, z')\}, \\ \forall(f, F) = (\lambda uz'\lambda z.\, f(uz)zz', \langle F, \mathtt{proj}_2\rangle) \end{cases}$$

for $\alpha \subseteq U \times X \times (Z \times Z')$ and $f : U \times Z \times Z' \to V,\ F : Y \times Z \times Z' \to X$. It just changes

$$\alpha(u, Fyzz', z, z') \supset \beta(fuzz', y, z, z')$$

into $\forall\alpha(u, Fyzz', z, z') \supset \forall\beta(\lambda z.f(uz)zz', y, z, z')$ which is further equivalent to

$$\alpha(uz, Fyzz', z, z') \supset \beta(f(uz)zz', y, z, z').$$

The functor $\exists$ is similarly defined.

The functors $\forall$ and $\exists$ are the right and left adjoints to the functor $\mathtt{proj}_2^{-1}$, respectively. The adjunction of $\mathtt{proj}_2^{-1}$ and $\forall$ is given by the bijection $(F, f) \mapsto (F, \lambda uz'\lambda z.\, fuzz')$ as seen by

$$\mathtt{proj}_2^{-1}(\alpha)(u, Fyzz', z, z') \supset \beta(fuzz', y, z, z')$$

and

$$\alpha(u, Fyzz', z') \supset \forall(\beta)((\lambda uz'\lambda z.\, fuzz')\, uz', y, z, z')$$

both of which are equivalent to

$$\alpha(u, Fyzz', z') \supset \beta(fuzz', y, z, z').$$

Similarly for the adjunction of $\exists$ and $\mathtt{proj}_2^{-1}$ due to the symmetry.

The Beck-Chevalley condition for $\forall$ is $f^{-1} \circ \forall(\alpha) \cong \forall \circ (\mathtt{id}_Z \times f)^{-1}(\alpha)$ which holds immediately, and similarly for $\exists$. Hence we have the following.

3.13. Theorem. *The indexed category $Z \mapsto \mathbb{GC}(Z)$ is a hyperdoctrine with $\mathbb{C} = \mathbb{SET}$.*

4. The double negation and ? ! translations

4.1. Gödel-Gentzen double-negation translation. Formulas of classical logic can be translated either into intuitionistic logic by $\neg\neg$-translation or into classical linear logic by ?!-translation. We hence have the situation depicted as below.

$$\begin{tikzcd}
\text{Classical Logic} \arrow[r, "?!"] \arrow[d, "\neg\neg"'] & \text{Classical Linear Logic} \arrow[d, "(\,)^L"] \\
\text{Intuitionistic Logic} \arrow[r, "(\,)^D"'] & \text{Higher-type } \exists\forall\text{-formulas}
\end{tikzcd}$$

where we regard linear logic as a subsystem of affine logic. We can in fact modify the $\neg\neg$- and ?!-translations to logically equivalent ones and establish that the modified ?!-translation and ϕ^L give the essentially equivalent $\exists\forall$-formulas, in the sense we will define, as the modified $\neg\neg$-translation and ϕ^D.

We start with the Gödel-Gentzen double-negation translation of classical logic to intuitionistic logic. It is defined inductively [Avigad and Feferman 1998] by

1. $\phi^N \equiv \neg\neg\phi$ for ϕ atomic,
2. $(\phi \wedge \psi)^N \equiv \phi^N \wedge \psi^N$,
3. $(\phi \vee \psi)^N \equiv \neg(\neg\phi^N \wedge \neg\psi^N)$,
4. $(\phi \supset \psi)^N \equiv \phi^N \supset \psi^N$,
5. $(\neg\phi)^N \equiv \neg\phi^N$,
6. $(\forall x\phi)^N \equiv \forall x\phi^N$,
7. $(\exists x\phi)^N \equiv \neg\forall x\neg\phi^N$

where the logical symbols on the left of $\equiv$ are classical while those on the right are intuitionistic.

This translation is known as the double-negation translation (a.k.a. $\neg\neg$-translation). In fact ϕ^N is equivalent to $\neg\neg\phi^N$ in intuitionistic logic. We can hence safely modify the Gödel-Gentzen double-negation translation to the following.

1. $\phi^{\mathcal{N}} \equiv \neg\neg\phi$ for ϕ atomic,
2. $(\phi \wedge \psi)^{\mathcal{N}} \equiv \neg\neg(\phi^{\mathcal{N}} \wedge \psi^{\mathcal{N}})$,
3. $(\phi \vee \psi)^{\mathcal{N}} \equiv \neg(\neg\phi^{\mathcal{N}} \wedge \neg\psi^{\mathcal{N}})$,
4. $(\neg\phi)^{\mathcal{N}} \equiv \neg\phi^{\mathcal{N}}$,
5. $(\forall x\phi)^{\mathcal{N}} \equiv \neg\neg\forall x\phi^{\mathcal{N}}$,
6. $(\exists x\phi)^{\mathcal{N}} \equiv \neg\forall x\neg\phi^{\mathcal{N}}$

where we no longer use $\supset$ as a primitive logical connective in classical logic.

4.2. PROPOSITION. *$\phi^{\mathcal{N}}$ is equivalent to ϕ^N in intuitionistic logic.*

PROOF. The proof is by induction on ϕ. It suffices to check the cases of $\phi \wedge \psi$ and $\forall x\phi$.

$(\phi \wedge \psi)^{\mathcal{N}} \equiv \neg\neg(\phi^{\mathcal{N}} \wedge \psi^{\mathcal{N}})$ is equivalent to $\neg\neg(\phi^N \wedge \psi^N) \equiv \neg\neg(\phi \wedge \psi)^N$ by the inductive hypothesis, which is in turn equivalent to $(\phi \wedge \psi)^N$.

$(\forall x\phi)^{\mathcal{N}} \equiv \neg\neg\forall x\phi^{\mathcal{N}}$ is equivalent to $\neg\neg\forall x\phi^N \equiv \neg\neg(\forall x\phi)^N$ by the inductive hypothesis which is in turn equivalent to $(\forall x\phi)^N$. ∎

4.3. GIRARD'S ? ! TRANSLATION. The Gödel-Gentzen double-negation translation can be understood as the interior-closure operation on open sets. In the phase semantics of linear logic the exponentials ? and ! are the closure and interior operations on facts. Girard extends this analogy to the translation of classical logic to linear logic [Girard 1987] given by

1. $\phi^\dagger \equiv\ ?\,!\phi$ for ϕ atomic,
2. $(\phi \vee \psi)^\dagger \equiv \phi^\dagger \parr \psi^\dagger$,
3. $(\neg\phi)^\dagger \equiv\ ?(\phi^{\dagger\perp})$,
4. $(\forall x\phi)^\dagger \equiv\ ?\,!\forall x\phi^\dagger$,
5. $(\phi \wedge \psi)^\dagger \equiv (\neg(\neg\phi \vee \neg\psi))^\dagger \equiv\ ?(!\phi^\dagger \otimes !\psi^\dagger)$,
6. $(\exists x\phi)^\dagger \equiv (\neg\forall x\neg\phi)^\dagger \equiv\ ?\,!\,?\exists x\,!(\phi^\dagger)$.

We call this translation ?!-translation. In fact $\phi^\dagger$ is equivalent to $?\,!\phi^\dagger$ in linear logic as seen from the following equivalences.

- $?\,!\phi$ and $?\,!\,?\,!\phi$,
- $?\,!\phi \parr ?\,!\psi$ and $?\,!(?\,!\phi \parr ?\,!\psi)$,
- $?(!\phi \otimes !\psi)$ and $?\,!\,?(!\phi \otimes !\psi)$.

Furthermore we can simplify the translation of $\exists x\phi$ to $?\exists x\,!\phi$ which is equivalent to $?\,!\,?\exists x\,!\phi$.

We can then safely modify Girard's ?!-translation to the following.

1. $\phi^+ =\ ?\,!\phi$ for ϕ atomic,
2. $(\phi \vee \psi)^+ =\ ?\,!\phi^+ \parr ?\,!\psi^+$,
3. $(\neg\phi)^+ =\ ?(\phi^{+\perp})$,
4. $(\forall x\phi)^+ =\ ?\,!\forall x\phi^+$,
5. $(\phi \wedge \psi)^+ =\ ?(!\phi^+ \otimes !\psi^+)$,
6. $(\exists x\phi)^+ =\ ?\exists x\,!\phi^+$.

4.4. PROPOSITION. *ϕ^+ is equivalent to $\phi^\dagger$ in linear logic.*

PROOF. The proof is by induction on ϕ. It suffices to check the disjunction. $(\phi \vee \psi)^\dagger = \phi^\dagger \invamp \psi^\dagger$ is equivalent to $?\,!\,\phi^\dagger \invamp\, ?\,!\,\psi^\dagger$, which is equivalent to $?\,!\,\phi^+ \invamp\, ?\,!\,\psi^+ = (\phi \vee \psi)^+$ by the inductive hypothesis. ∎

4.5. RELATING THE TWO TRANSLATIONS. Let us consider the set VT of terms which are constructed from variables and application only. VT is defined as the smallest set satisfying

- $u \in VT$ for any variable u,
- if $s \in VT$ and $t \in VT$ then $st \in VT$.

Any term s in VT has the form

$$u_s t_1 t_2 \dots t_n$$

with a variable u_s. We call u_s the head variable of s as usual.

Let us then consider the set Θ of pairs (p, q) of closed terms

$$\begin{cases} p : (\tau_1 \to \tau_2 \to \cdots \to \tau_n \to \mathsf{0}) \to (\tau'_{\pi(1)} \to \tau'_{\pi(2)} \to \cdots \to \tau'_{\pi(n)} \to \mathsf{0}) \\ q : (\tau'_{\pi(1)} \to \tau'_{\pi(2)} \to \cdots \to \tau'_{\pi(n)} \to \mathsf{0}) \to (\tau_1 \to \tau_2 \to \cdots \to \tau_n \to \mathsf{0}) \end{cases}$$

defined inductively on the construction of types $\tau_1 \to \cdots \to \tau_n \to \mathsf{0}$ so that

- $(\lambda x.\, x, \lambda x.\, x) \in \Theta$ for variables x of the base type,
- if $(p_i, q_i) \in \Theta$ with $p_i : \tau_i \to \tau'_i$, $q_i : \tau'_i \to \tau_i$ for $1 \le i \le n$ and π is a permutation on $\{1, 2, \dots, n\}$ then $(p, q) \in \Theta$ for

$$\begin{cases} p \equiv \lambda u x_{\pi(1)} x_{\pi(2)} \dots x_{\pi(n)}.\, u(q_1 x_1)(q_2 x_2) \dots (q_n x_n) \\ q \equiv \lambda u x_1 x_2 \dots x_n.\, u(p_{\pi(1)} x_{\pi(1)})(p_{\pi(2)} x_{\pi(2)}) \dots (p_{\pi(n)} x_{\pi(n)}). \end{cases}$$

The intended interpretation of p in a Cartesian closed category $\mathbb{C}$ is an isomorphism

$$(A_1 \to A_2 \to \cdots \to A_n \to D) \to (A'_{\pi(1)} \to A'_{\pi(2)} \to \cdots \to A'_{\pi(n)} \to D)$$

generated from the family of canonical isomorphisms

$$A \times B \to B \times A, \qquad (A \times B) \times C \to A \times (B \times C)$$

by composition, the adjunction $\mathbb{C}(A \times B, C) \cong \mathbb{C}(B, A \to C)$, the functor $(\_) \times (\_)$ and the functor $(\_) \to D$ for a fixed object D. The interpretation of q is its inverse.

4.6. LEMMA. *Let $(p, q) \in \Theta$. Then $q(pu) = u$ and $p(qu) = u$.*

PROOF. The proof is by induction on types. One of the inductive case is shown as follows.

$$\begin{aligned} q(pu) &= \lambda x_1 x_2 \dots x_n.\, u(q_1(p_1 x_1))(q_2(p_2 x_2)) \dots (q_n(p_n x_n)) \\ &= \lambda x_1 x_2 \dots x_n.\, u x_1 x_2 \dots x_n = u. \end{aligned}$$

∎

Let θ be a function from the pairs $u^\sigma = (u, \sigma)$ of a variable x and a type σ to Θ such that $p_{(u,\sigma)} : \sigma \to \tau$ and $p_{(u,\tau)} = q_{(u,\sigma)}$ where we write $(p_{(u,\sigma)}, q_{(u,\sigma)})$ for $\theta(u, \sigma)$. We suppress the reference to types as much as possible and often simply write $\theta(u)$ and so on. We also write π_u for the permutation π such that

$$p_u : (\tau_1 \to \tau_2 \to \cdots \to \tau_n \to \mathsf{0}) \to (\tau'_{\pi_u(1)} \to \tau'_{\pi_u(2)} \to \cdots \to \tau'_{\pi_u(n)} \to \mathsf{0})$$

while π_u is the empty function if u has the base type. Furthermore we write p_i^u and q_i^u for p_i and q_i used in the inductive definitions of p_u and q_u, and define

$$\begin{cases} p_s \equiv \lambda v x_{\pi_u(k+1)} x_{\pi_u(k+2)} \ldots x_{\pi_u(n)}.\, v(q_{k+1}^u x_{k+1})(q_{k+2}^u x_{k+2}) \ldots (q_n^u x_n) \\ q_s \equiv \lambda v x_{k+1} x_{k+2} \ldots x_n.\, v(p_{\pi_u(k+1)}^u x_{\pi_u(k+1)})(p_{\pi_u(k+2)}^u x_{\pi_u(k+2)}) \ldots (p_{\pi_u(n)}^u x_{\pi_u(n)}) \end{cases}$$

for $s \equiv u s_1 s_2 \ldots s_k$ when the restriction of π_u to $\{k+1, k+2, \ldots, n\}$ is a permutation. They are intended as isomorphisms for the type of s.

4.7. LEMMA. *$\pi_{(u,\sigma)} = \pi_{(u,\tau)}^{-1}$ when $p_{(u,\sigma)} : \sigma \to \tau$.*

PROOF. Suppose otherwise. We then have $p_i^{(u,\sigma)} y_i = q_j^{(u,\tau)} y_j$ for $i \neq j$. Substitute $\lambda w.0$ and $\lambda w.1$ for y_i and y_j, respectively, and derive the contradiction. ∎

4.8. LEMMA. *$q_{(u,\tau)} = p_{(u,\sigma)}$ and $p_i^{(u,\sigma)} = q_i^{(u,\tau)}$ when $p_{(u,\sigma)} : \sigma \to \tau$.*

PROOF. For the first equation we have

$$q_{(u,\tau)} x = p_{(u,\sigma)}(q_{(u,\sigma)}(q_{(u,\tau)} x)) = p_{(u,\sigma)}(p_{(u,\tau)}(q_{(u,\tau)} x)) = p_{(u,\sigma)} x.$$

For the second equation let $s = \lambda y_1 \ldots y_n.\, y_i z_1 \ldots z_k$. Then

$$p_{(u,\tau)}(q_{(u,\tau)} s) = p_{(u,\tau)}(p_{(u,\sigma)} s) = \lambda y_1 y_2 \ldots y_n.\, q_i^{(u,\sigma)}(q_i^{(u,\tau)} y_i) z_1 \ldots z_k = s.$$

We hence have $q_i^{(u,\sigma)}(q_i^{(u,\tau)} y_i) z_1 \ldots z_k = y_i z_1 \ldots z_k$ and

$$p_i^{(u,\sigma)} y_i = p_i^{(u,\sigma)}(q_i^{(u,\sigma)}(q_i^{(u,\tau)} y_i)) = q_i^{(u,\tau)} y_i.$$

∎

4.9. DEFINITION. *For a given θ we define the relation $s \sim_\theta s'$ on VT as the smallest relation satisfying*

- *$u^\sigma \sim_\theta u^\tau$ for any variable u if $p_{(u,\sigma)} : \sigma \to \tau$,*
- *if $u^\sigma \sim_\theta u^\tau$ and $s_i \sim_\theta t_i$ for $1 \le i \le k \le n$ so that*
  - *the restriction of π_u on $\{1, 2, \ldots, n\}$ to $\{1, 2, \ldots, k\}$ is a permutation,*
  - *$p_i^u = p_{s_i}$ and $q_i^u = q_{s_i}$ for $1 \le i \le k$*

  then $u^\sigma s_1 s_2 \ldots s_n \sim_\theta u^\tau t_{\pi_u(1)} t_{\pi_u(2)} \ldots t_{\pi_u(n)}$.

If there is θ such that $s \sim_\theta t$ we simply say s is essentially equivalent to t.

4.10. LEMMA. *The relation $s \sim_\theta s'$ is symmetric.*

PROOF. By induction on s using Lemma 4.8 ∎

4.11. LEMMA. *Suppose the list x of variables in s is a permutation of y in t and x does not include the variable u. If s is essentially equivalent to t then $s[ux/u]$ is essentially equivalent to $t[uy/u]$.*

PROOF. Suppose $s \sim_{\theta'} t$ and $y_1, y_2, \ldots, y_k$ is $x_{\pi(1)}, x_{\pi(2)}, \ldots, x_{\pi(k)}$. We write (p'_u, q'_u) for $\theta'(u)$ and so on. Let $\theta(u)$ be such that

$$\pi_u(i) = \begin{cases} \pi(i) & \text{for} \quad 1 \le i \le k \\ \pi'_u(i-k)+k & \text{for} \quad k < i \le n, \end{cases} \qquad p^u_i = \begin{cases} p'_{x_i} & \text{for} \quad 1 \le i \le k \\ p'^u_{i-k} & \text{for} \quad k < i \le n \end{cases}$$

and similarly for q^u_i. Let $\theta(v) = \theta'(v)$ for $v \not\equiv u$. ∎

4.12. THEOREM. *Suppose $s \sim_\theta t$ and the head variable of s is u. Suppose further $x_{k+1}, x_{k+2}, \ldots, x_n$ is the list of fresh variables so that $s x_{k+1} x_{k+2} \ldots x_n$ is a term of $\mathsf{0}$. Then*

$$s x_{k+1} x_{k+2} \ldots x_n = t[pv/v] r_{\pi_u(k+1)} r_{\pi_u(k+2)} \ldots r_{\pi_u(n)}$$

where $r_i \equiv p^u_i x_i$ for $k+1 \le i \le n$ and pv is the list $p_{v_1} v_1, \ldots, p_{v_j} v_j$ with v including all the variables of s.

PROOF. The proof is by induction on the construction of s. If s is a variable u then

$$pu r_{\pi_u(1)} r_{\pi_u(2)} \ldots r_{\pi_u(n)} = u(q_1(p_1 x_1)) \ldots (q_n(p_n x_n)) = u x_1 \ldots x_n.$$

Suppose $s \equiv u s_1 \ldots s_k$ and $t \equiv u t_{\pi_u(1)} \ldots t_{\pi_u(k)}$. Suppose further the head variable of s_i is w. By the inductive hypothesis we have

$$s_i y_{l+1} y_{l+2} \ldots y_m = t_i[pv/v] r_{\pi_w(l+1)} r_{\pi_w(l+2)} \ldots r_{\pi_w(m)}$$

where $r_j \equiv p^w_j y_j$ for $l+1 \le j \le m$. We then have

$$\begin{aligned} q^u_i(t_i[pv/v]) y_{l+1} \ldots y_m &= q_{s_i}(t_i[pv/v]) y_{l+1} \ldots y_m \\ &= t_i[pv/v](p^w_{\pi_w(l+1)} y_{\pi_w(l+1)}) \ldots (p^w_{\pi_w(m)} y_{\pi_w(m)}) \\ &= s_i y_{l+1} \ldots y_m \end{aligned}$$

so that $q^u_i(t_i[pv/v]) = s_i$. Hence

$$\begin{aligned} & t[pv/v] r_{\pi_u(k+1)} \ldots r_{\pi_u(n)} \\ &\quad = pu(t_{\pi_u(1)}[pv/v]) \ldots (t_{\pi_u(k)}[pv/v]) r_{\pi_u(k+1)} \ldots r_{\pi_u(n)} \\ &\quad = u(q^u_1(t_1[pv/v])) \ldots (q^u_k(t_k[pv/v]))(q^u_{k+1}(p^u_{k+1} x_{k+1})) \ldots (q^u_n(p^u_n x_n)) \\ &\quad = u s_1 \ldots s_k x_{k+1} \ldots x_n. \end{aligned}$$

∎

4.13. Corollary. *Use the same assumptions as Theorem 4.12. Suppose further that y and w are disjoint lists of variables such that every variable in v is either in y or w. Then*

$$s[qy/y]x_{k+1}x_{k+2}\dots x_n = t[pw/w]r_{\pi_u(k+1)}r_{\pi_u(k+2)}\dots r_{\pi_u(n)}.$$

Proof.

$$\begin{aligned} s[qy/y]x_{k+1}\dots x_n &= (t[pv/v])\,[qy/y]r_{\pi_u(k+1)}\dots r_{\pi_u(n)} \\ &= t[pw/w,\, p_{y_1}(q_{y_1}y_1)/y_1,\dots,p_{y_m}(q_{y_m}y_m)/y_m]r_{\pi_u(k+1)}\dots r_{\pi_u(n)} \\ &= t[pv/v]r_{\pi_u(k+1)}\dots r_{\pi_u(n)} \end{aligned}$$

∎

Let u and x be the lists of variables of type $\mathtt{0}$ with the length m and n, respectively. Furthermore let U, U' be the lists of terms of $\mathtt{0}$ in VT with the length m and X, X' be the lists of terms of $\mathtt{0}$ in VT with the length n.

4.14. Definition. *α and α' are essentially equivalent under θ, denoted $\alpha \sim_\theta \alpha'$, if*

- *there are quantifier-free formulas $\beta(u,x,z)$ and $\beta'(u,x,z)$ such that $\beta \leftrightarrow \beta'$ holds, which contain exactly the same free variables of type $\mathtt{0}$ and no others,*
- *$U_i \sim_\theta U'_i$ and $X_j \sim_\theta X'_j$ for each $1 \le i \le m$ and $1 \le j \le n$,*
- *$\alpha \equiv \beta(U,X,z)$ and $\alpha' \equiv \beta'(U',X',z)$.*

If $\alpha \sim_\theta \alpha'$ for some θ we simply say α is essentially equivalent to α', denoted $\alpha \sim \alpha'$. The terms U, U', X and X' are called primary terms.

4.15. Theorem. *For any formula ϕ of first-order classical logic $\left(\phi^{\mathcal{N}}\right)_D$ and $(\phi^+)_L$ are essentially equivalent.*

The proof is by induction on the construction of formulas and we use Lemma 4.11 to conclude $\alpha[ux/u] \sim \beta[uy/u]$ from $\alpha \sim \beta$ when y is a permutation of x.

Nothing is to be proved for the atomic case. Assuming $\phi_D(u,x) \sim \psi_L(u,x)$ we have

$$\begin{aligned} (\neg\phi)_D &\equiv \neg\phi_D(u,xu) \sim \neg\psi_L(u,xu) \equiv \left(?\,\phi^\perp\right)_L, \\ (\neg\neg\phi)_D &\equiv \neg\neg\phi_D(ux,x(ux)) \sim \psi_L(ux,x(ux)) \equiv (\,?\,!\,\psi)_L \equiv \left(?\,(\,?\,\psi^\perp)^\perp\right)_L. \end{aligned}$$

The cases of $\neg\phi$ and $\forall x\phi$ follow from them. The rest of the cases are treated by the following step-by-step calculations.

$(\phi \wedge \psi)^{\mathcal{N}}$ and $(\phi \wedge \psi)^{+}$. $\left((\phi \wedge \psi)^{\mathcal{N}}\right)_D$ is computed as follows.

$$
\begin{array}{ccc}
\phi^{\mathcal{N}} & & \psi^{\mathcal{N}} \\
& \diagdown \quad \diagup & \\
& \phi^{\mathcal{N}} \wedge \psi^{\mathcal{N}} & \\
& | & \\
& \neg(\phi^{\mathcal{N}} \wedge \psi^{\mathcal{N}}) & \\
& | & \\
& \neg\neg(\phi^{\mathcal{N}} \wedge \psi^{\mathcal{N}}) &
\end{array}
\qquad
\begin{array}{ccc}
\left(\phi^{\mathcal{N}}\right)_D(u,x) & & \left(\psi^{\mathcal{N}}\right)_D(v,y) \\
& \diagdown \quad \diagup & \\
& \left(\phi^{\mathcal{N}}\right)_D(u,x) \wedge \left(\psi^{\mathcal{N}}\right)_D(v,y) & \\
& | & \\
& \neg(\left(\phi^{\mathcal{N}}\right)_D(u,xuv) \wedge \left(\psi^{\mathcal{N}}\right)_D(v,yuv)) & \\
& | & \\
& \neg\neg(\left(\phi^{\mathcal{N}}\right)_D(uxy, x(uxy)(vxy)) \wedge \left(\psi^{\mathcal{N}}\right)_D(vxy, y(uxy)(vxy))) &
\end{array}
$$

$((\phi \wedge \psi)^{+})_L$ is as follows. We only need to permute u and v in xuv in the third step.

$$
\begin{array}{ccc}
\phi^{+} & & \psi^{+} \\
| & & | \\
!\,\phi^{+} & & !\,\psi^{+} \\
& \diagdown \quad \diagup & \\
& !\,\phi^{+} \otimes\, !\,\psi^{+} & \\
& | & \\
& ?\,(\,!\,\phi^{+} \otimes\, !\,\psi^{+}) &
\end{array}
\qquad
\begin{array}{ccc}
(\phi^{+})_L(u,x) & & (\psi^{+})_L(v,y) \\
| & & | \\
(\phi^{+})_L(u,xu) & & (\psi^{+})_L(v,yv) \\
& \diagdown \quad \diagup & \\
& (\phi^{+})_L(u,xvu) \wedge (\psi^{+})_L(v,yuv) & \\
& | & \\
& (\phi^{+})_L(uxy, x(vxy)(uxy)) \wedge (\psi^{+})_L(vxy, y(uxy)(vxy)) &
\end{array}
$$

$(\phi \vee \psi)^{\mathcal{N}}$ and $(\phi \vee \psi)^{+}$. $\left((\phi \vee \psi)^{\mathcal{N}}\right)_D$ is as follows.

$$
\begin{array}{ccc}
\phi^{\mathcal{N}} & & \psi^{\mathcal{N}} \\
| & & | \\
\neg\phi^{\mathcal{N}} & & \neg\psi^{\mathcal{N}} \\
& \diagdown \quad \diagup & \\
& \neg\phi^{\mathcal{N}} \wedge \neg\psi^{\mathcal{N}} & \\
& | & \\
& \neg\left(\neg\phi^{\mathcal{N}} \wedge \neg\psi^{\mathcal{N}}\right) &
\end{array}
\qquad
\begin{array}{ccc}
\left(\phi^{\mathcal{N}}\right)_D(u,x) & & \left(\psi^{\mathcal{N}}\right)_D(v,y) \\
| & & | \\
\neg\left(\phi^{\mathcal{N}}\right)_D(u,xu) & & \neg\left(\psi^{\mathcal{N}}\right)_D(v,yv) \\
& \diagdown \quad \diagup & \\
& \neg\left(\phi^{\mathcal{N}}\right)_D(u,xu) \wedge \neg\left(\psi^{\mathcal{N}}\right)_D(v,yv) & \\
& | & \\
& \neg(\neg\left(\phi^{\mathcal{N}}\right)_D(uxy, x(uxy)) \wedge \neg\left(\psi^{\mathcal{N}}\right)_D(vxy, y(vxy))) &
\end{array}
$$

$((\phi \vee \psi)^{+})_L$ is as follows. We only need to permute x and y in uxy in the last step.

$$
\begin{array}{ccc}
\phi^{+} & & \psi^{+} \\
| & & | \\
!\,\phi^{+} & & !\,\psi^{+} \\
| & & | \\
?\,!\,\phi^{+} & & ?\,!\,\psi^{+} \\
& \diagdown \quad \diagup & \\
& ?\,!\,\phi^{+} \parr\, ?\,!\,\psi^{+} &
\end{array}
\qquad
\begin{array}{ccc}
(\phi^{+})_L(u,x) & & (\psi^{+})_L(v,y) \\
| & & | \\
(\phi^{+})_L(u,xu) & & (\psi^{+})_L(v,yv) \\
| & & | \\
(\phi^{+})_L(ux, x(ux)) & & (\psi^{+})_L(vy, y(vy)) \\
& \diagdown \quad \diagup & \\
& (\phi^{+})_L(uyx, x(uyx)) \vee (\psi^{+})_L(vxy, y(vxy)) &
\end{array}
$$

$(\exists y\, \phi)^{\mathcal{N}}$ **and** $(\exists y\, \phi)^{+}$**.** $\left((\exists y\, \phi)^{\mathcal{N}}\right)_D$ is as follows.

$$\begin{array}{c} \phi^{\mathcal{N}} \\ | \\ \neg\phi^{\mathcal{N}} \\ | \\ \forall y\, \neg\phi^{\mathcal{N}} \\ | \\ \neg\forall y\, \neg\phi^{\mathcal{N}} \end{array} \qquad\qquad \begin{array}{c} \left(\phi^{\mathcal{N}}\right)_D(u, x, y) \\ | \\ \neg\left(\phi^{\mathcal{N}}\right)_D(u, xu, y) \\ | \\ \neg\left(\phi^{\mathcal{N}}\right)_D(u, xyu, y) \\ | \\ \neg\neg\left(\phi^{\mathcal{N}}\right)_D(ux, x(yx)(ux), yx) \end{array}$$

$\left((\exists y\, \phi)^{+}\right)_L$ is as follows. There is no need for permutation.

$$\begin{array}{c} \phi^{+} \\ | \\ !\, \phi^{+} \\ | \\ \exists y\ !\, \phi^{+} \\ | \\ ?\, \exists y\ !\, \phi^{+} \end{array} \qquad\qquad \begin{array}{c} (\phi^{+})_L(u, x, y) \\ | \\ (\phi^{+})_L(u, xu, y) \\ | \\ (\phi^{+})_L(u, xyu, y) \\ | \\ (\phi^{+})_L(ux, x(yx)(ux), yx) \end{array}$$

4.16. Corollary. *$\left(\phi^{\mathcal{N}}\right)_D(u, x)$ has witnesses if and only if $(\phi^{+})_L(u, x)$ has witnesses.*

Proof. Let $\left(\phi^{\mathcal{N}}\right)_D(u, x) \equiv \alpha(U, X)$ and $(\phi^{+})_L(u, x) \equiv \beta(U', X')$. Assuming that $\left(\phi^{\mathcal{N}}\right)_D(\boldsymbol{u}, x) \equiv \alpha(U[\boldsymbol{u}/u], X[\boldsymbol{u}/u])$ holds we have

$$\alpha(U[\boldsymbol{u}/u,\ qx/x], X[\boldsymbol{u}/u,\ qx/x])$$

by substitution for negative variables. Since primary terms are of type $\mathtt{0}$ we have

$$\alpha(U[qx/x], X[qx/x]) \leftrightarrow \beta(U'[pu/u], X'[pu/u])$$

by Corollary 4.13 and Theorem 4.15. We then have

$$\alpha(U[\boldsymbol{u}/u,\ qx/x], X[\boldsymbol{u}/u,\ qx/x]) \leftrightarrow \beta(U'[p\boldsymbol{u}/u], X'[p\boldsymbol{u}/u])$$

by substitution for positive variables. It follows from them that

$$\left(\phi^{+}\right)_L(p\boldsymbol{u}, x) \equiv \beta(U'[p\boldsymbol{u}/u], X'[p\boldsymbol{u}/u])$$

holds. The other direction is entirely similar. ∎

4.17. Shoenfield's version of Dialectica translation. In his textbook *Mathematical Logic*, Shoenfield gave a consistency proof of *PA*, using a version of Dialectica interpretation which interprets classical formulas directly [Shoenfield 1967]. His version of D-translation is the translation of a classical formula ϕ into a classical higher-type formula ϕ^S of the form

$$\forall x \exists u\, \phi_S(u, x).$$

which is defined inductively by

- $\phi^S \equiv \phi_S \equiv \phi$ for atomic ϕ,
- $(\neg\phi)^S \equiv \forall u \exists x\, \neg\phi_S(ux, x)$,
- $(\phi \vee \psi)^S \equiv \forall xy \exists uv\ (\phi_S(u, x) \vee \psi_S(v, y))$,
- $(\forall z\, \phi)^S \equiv \forall zx \exists u\, \phi_S(u, x, z)$.

$\phi \wedge \psi$ and $\exists z\, \phi$ are defined as $\neg\,(\neg\phi \vee \neg\psi)$ and $\neg\forall z\, \neg\phi$, respectively. They do not appear in Shoenfield's consistency proof at all. If we compute their D-translations according to the definitions, however, they become the following.

- $(\phi \wedge \psi)^S \equiv \forall xy \exists uv\, \neg\,(\neg\phi_S(u(xuv), xuv) \vee \neg\psi_S(v(yuv), yuv))$.
- $(\exists z\, \phi)^S \equiv \forall x \exists zu\, \phi_S(u(xuz), xuz, z, z')$.

Shoenfield then showed that if ϕ is provable in *PA* then there exists the list $\boldsymbol{u}$ of witness terms such that $\phi_S(\boldsymbol{u}, x)$ holds. The witnesses are, however, allowed to contain free variables from x. Hence we can instead consider the Skolemization

$$\phi^{SS} \equiv \exists u \forall x\, \phi_{SS}(u, x) \equiv \exists u \forall x\, \phi_S(ux, x)$$

of ϕ^S and obtain witnesses with no free variables from x.

Let us regard our D-interpretations of linear logic operators as the operations on quantifier-free formulas with designated positive and negative variables. We then have the following observation.

- $\phi_{SS} \equiv \phi_S(ux, x) \equiv\ ?\,\phi_S(u, x)$.
- If ϕ is atomic, then $\phi_{SS} \equiv \phi \equiv\ ?\,!\,\phi$ since ϕ has no bound variables.
- $(\neg\phi)_{SS} \equiv \neg\phi_S(u(xu), xu) \equiv\ ?\,(\,?\,\phi_S(u, x))^{\perp} \equiv\ ?\,(\phi_{SS})^{\perp}$.
- $(\phi \vee \psi)_{SS} \equiv \phi_S(uxy, x) \vee \psi_S(vxy, y) \sim\ ?\,\phi_S(u, x) \parr\ ?\,\psi_S(v, y) \sim \phi_{SS} \parr \psi_{SS}$.
- $(\phi \wedge \psi)_{SS} \equiv$
  $\neg\,(\neg\phi_S(uxy(x(uxy)(vxy)), x(uxy)(vxy)) \vee \neg\psi_S(vxy(y(uxy)(vxy)), y(uxy)(vxy)))$
  $\sim\ ?\,(\,!\,?\,\phi_S(u, x) \otimes\ !\,?\,\psi_S(v, y)) \equiv\ ?\,(\,!\,\phi_{SS} \otimes\ !\,\psi_{SS})$.

- $(\forall z\,\phi)_{SS} \equiv \phi_S(uxz, x, z) \equiv \forall z \;?\,\phi_S(u, x, z) \equiv \forall z \phi_{SS}$.
- $(\exists z\,\phi)_{SS} \equiv \phi_S(ux(x(ux)(zx)), x(ux)(zx), zx) \sim \;?\,\exists z \;!\;?\,\phi_S(u, x, z) \equiv \;?\,\exists z \;!\,\phi_{SS}$.

ϕ^{SS} thus corresponds almost exactly to the composition of ?!-translation and our D-translation, where we use $\phi^{\dagger}$ for the disjunction and ϕ^{+} for the existential quantifier. The only exception is the universal quantification. $\left((\forall z\phi)^{+}\right)^{L}$ corresponds to $(\neg\neg\forall z\phi)^{SS}$ rather than $(\forall z\phi)^{SS}$.

Hyland analyzed $\phi^{S} \equiv \forall x \exists u\, \phi_S(u, x)$ by converting it into $\phi_S(ux, x) \supset \bot$, which can be regarded as the object $R^{\bar{\phi}}$ in the categorical setting, where $\bar{\phi}$ and R are objects $\phi_S(u, x)$ and $\bot$, respectively [Hyland 2002]. We can write $\phi_S(ux, x) \supset \bot$ as

$$?\,\neg\phi_S(ux, x) \equiv \;?\,(\phi_S(u, x))^{\perp}$$

and our description of ϕ^{SS} fits with Hyland's analysis. We note however that we are able to give a more direct description due to the presence of ?.

4.18. The translation $\phi^{\flat}$. It is interesting to think about why $\neg\neg$ is not necessary in $(\forall x\phi)^{SS}$. This is also the case in the combination of the original $\neg\neg$-translation ϕ^{N} and ϕ^{D}. Shoenfield's D-translation suggests the following translation $\phi^{\flat}$ of classical logic into classical linear logic.

1. $\phi^{\flat} \equiv \;?\,\phi$ for ϕ atomic.
2. $(\phi \vee \psi)^{\flat} \equiv \phi^{\flat} \parr \psi^{\flat}$.
3. $(\neg\phi)^{\flat} \equiv \;?\left(\phi^{\flat}\right)^{\perp}$.
4. $(\phi \wedge \psi)^{\flat} \equiv \;?\left(!\,\phi^{\flat} \otimes \,!\,\psi^{\flat}\right)$.
5. $(\forall x\phi)^{\flat} \equiv \forall x \phi^{\flat}$.
6. $(\exists x\phi)^{\flat} \equiv \;?\,\exists x \;!\,\phi^{\flat}$.

$\phi^{\flat}$ is logically equivalent to $?\,\phi^{\flat}$. Furthermore $\phi^{\flat}$ is sufficient to capture the provability of classical logic in the following sense.

4.19. Theorem. *If a sequent $\phi_1, \ldots, \phi_n \vdash \psi_1, \ldots, \psi_m$ is provable in classical logic, then the sequent*

$$!\,{\phi_1}^{\flat}, \ldots, !\,{\phi_n}^{\flat} \vdash {\psi_1}^{\flat}, \ldots, {\psi_n}^{\flat}$$

is provable in classical linear logic.

Proof. The proof is by induction on the construction of classical proof. Interesting cases are the left $\vee$ rule and the left $\forall$ rule, where we use the fact that

$$!\left(\phi^{\flat} \parr \psi^{\flat}\right) \vdash \;?\,!\,\phi^{\flat} \parr \;?\,!\,\psi^{\flat} \qquad \text{and} \qquad !\,\forall x \phi^{\flat} \vdash \forall x \;!\,\phi^{\flat}$$

are provable in classical linear logic. ∎

In classical affine logic $(\phi \wedge \psi)^\flat$ can be simplified to $?\left(\phi^\flat \otimes \psi^\flat\right)$ for the similar result, since $!\,?\left(\phi^\flat \otimes \psi^\flat\right) \vdash\ !\,\phi^\flat \otimes\ !\,\psi^\flat$ is then provable, using the left weakening on $\phi^\flat$ and $\psi^\flat$.

5. Conclusion

We have obtained a structurally symmetric reformulation of Gödel's Dialectica interpretation by way of a contraction-free logic. We would like to finish our work by making three additional comments.

First of all although our work is motivated mostly by a theoretical concern our interpretation may be useful for practical purposes such as program extraction. There have been attempts to refine a logical system so that one can extract optimized programs from proofs, for example Berger's uniform Heyting arithmetic [Berger 2005]. By keeping track of the use of contraction we can certainly curtail unnecessary dependencies and obtain a smaller witness term.

Secondly when we extend our interpretation to arithmetic we need to add to our logical system the axioms for equality and primitive recursive functions, and the rule of inference for induction. Every atomic formula $s = t$ is then contractible since we can derive

$$s = t \vdash s = t \otimes s = t$$

from the axioms

$$\vdash s = s, \qquad s = t,\, s = s \otimes s = s \vdash s = t \otimes s = t$$

where we use the two-sided sequent calculus for readability. Things are otherwise straightforward up to the induction rule, of which we can consider the four versions

$$\frac{\vdash \phi(0) \quad \phi(x) \vdash \phi(sx)}{\vdash \phi(t)} \qquad \frac{\vdash \phi(0) \quad !\,\phi(x) \vdash \phi(sx)}{\vdash \phi(t)}$$

$$\frac{\Gamma,\, !\,\Delta \vdash \phi(0) \quad !\,\Delta, \phi(x) \vdash \phi(sx)}{\Gamma,\, !\,\Delta \vdash \phi(t)} \qquad \frac{!\,\Gamma,\, !\,\Delta \vdash \phi(0) \quad !\,\Delta,\, !\,\phi(x) \vdash \phi(sx)}{!\,\Gamma,\, !\,\Delta \vdash \phi(t)}$$

where s is the successor and t is an arbitrary term. If we allow ϕ to be an arbitrary formula they are admissible from one another. The interesting thing would be, however, to investigate the Dialectica interpretations of weaker forms of induction, obtained by restricting ϕ to a certain class of formulas or restricting the exponentials as Girard's light linear logic [Girard 1998]. This would be analogous to the work of Cook and Urquhart on bounded arithmetic [Cook and Urquhart 1993].

Finally our Dialectica interpretation seems to be related to two types of games. One is the game semantics of propositional linear logic and PCF [Abramsky and Jagadeesan 1994, Abramsky et al. 2000]. The other is the Henkin-Hintikka game for first-order classical logic [Hintikka 1996]. Our study has been guided by the observation that the game semantics of PCF is immediately applicable to its subsystem T, and by the intuition that the Dialectica interpretation is a constructive version of the Henkin-Hintikka game.

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